

Interference Evaluation in ECMA-368 WPAN with DRP MAC Protocol

Holger Rosier

ComNets Research Group

RWTH Aachen University, Germany

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Motivation

- Wireless communication in residential home environments covers wide range of applications
- Common market trend from multimedia capable devices (DEVs) towards applications demanding for very high data rate transmission service: TV and video data transmission
- Outdated Wifi's mono-cluster approach replaced by multi-cluster home networks yielding intended high data throughput in limited area ^[1]
- ➔ Motivation to use UWB ECMA -368 for short range ad-hoc communication in residential home environment

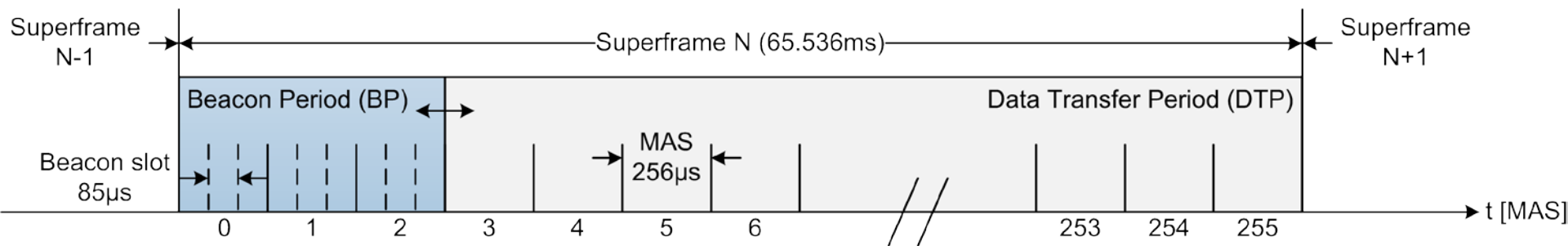
[1] P. Gandolfo. XtremeSpectrum – SG3a CFA Response. Online available: 02031r0P802-15_SGAP3-CFAReapponseAltPHY

Problem Description

- Widespread use of wireless communication systems operating in license-exempt frequency bands goes along densely populated network scenarios
- DEV applies reservation based DRP controlled medium access to transmit user data reliably carried in successively recurring Medium Access Slots (MASs)
- Spatial reuse of MAS channel resources in not fully meshed network scenarios causes interference to user data frame transmission
- ➔ Limits achievable cluster data throughput and coverage range

ECMA-368^[2]: Medium Access Control (MAC) Overview

- Fully distributed MAC:
 - DEVs unhide DEV located in the neighbourhood through short management frames broadcasted during Beacon Period (BP)
 - DEV negotiates MASs for user data frame transmission in advance to carry user data frames during Data Transfer Period (DTP)



- Management frames carried in Beacon Slots (BSs) are protected from interference caused by transmission in the two hop neighbourhood
- User data frames carried in MAS exclusively are protected from interference caused by DEVs data transmissions located in one hop Beacon Group (BG)

[2] ECMA-368: High Rate Ultra Wideband PHY and MAC Standard, 3rd Edition, December 2008.

ECMA-368: Medium Access Control Beacon Period (BP)

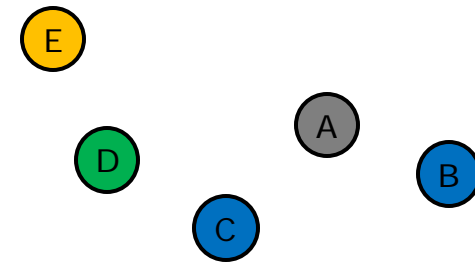
- Beacons are broadcasted in recurring BSs
- Several Information Elements (IEs) compose beacon
- DEV are aware of DEVs located in one and two hop neighbourhood, defers from trying to reserve respective BSs
- Beacon Decoding Range is rather limited by noise than interference caused from data frame transmission

r_{bg} Beacon Decoding Range



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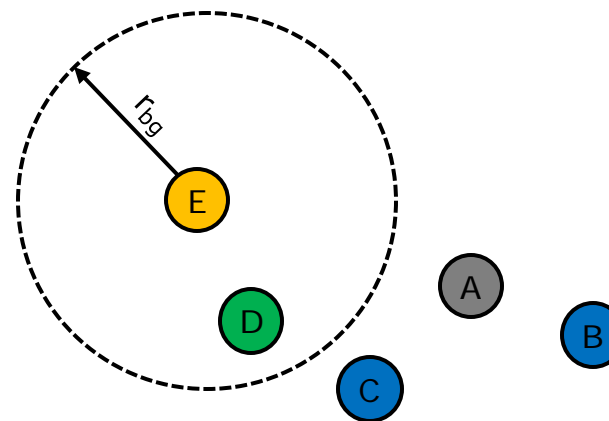


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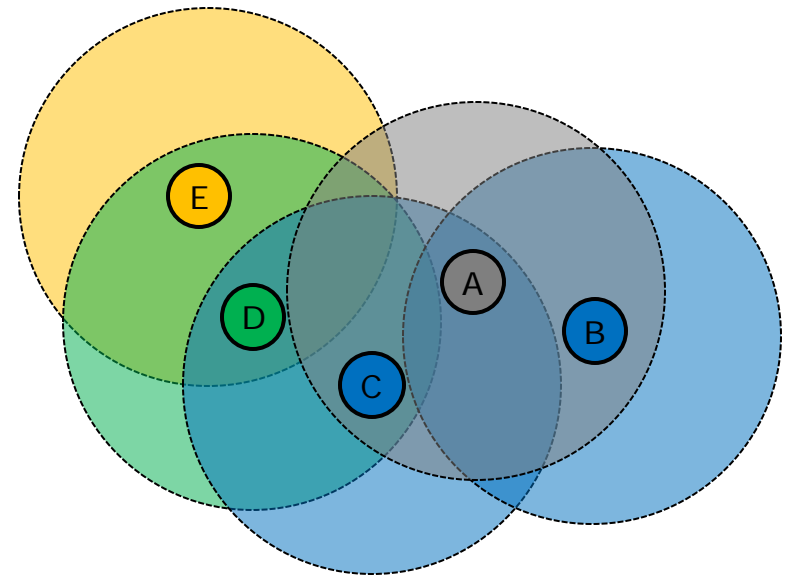


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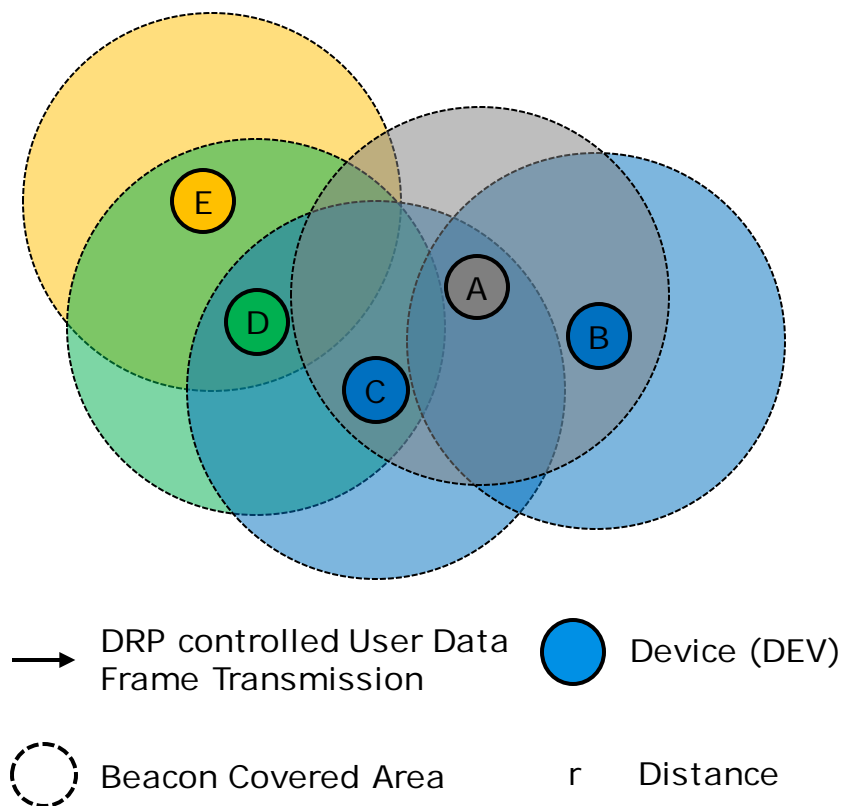


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ECMA-368: Medium Access Control Data Transfer Period (DTP)

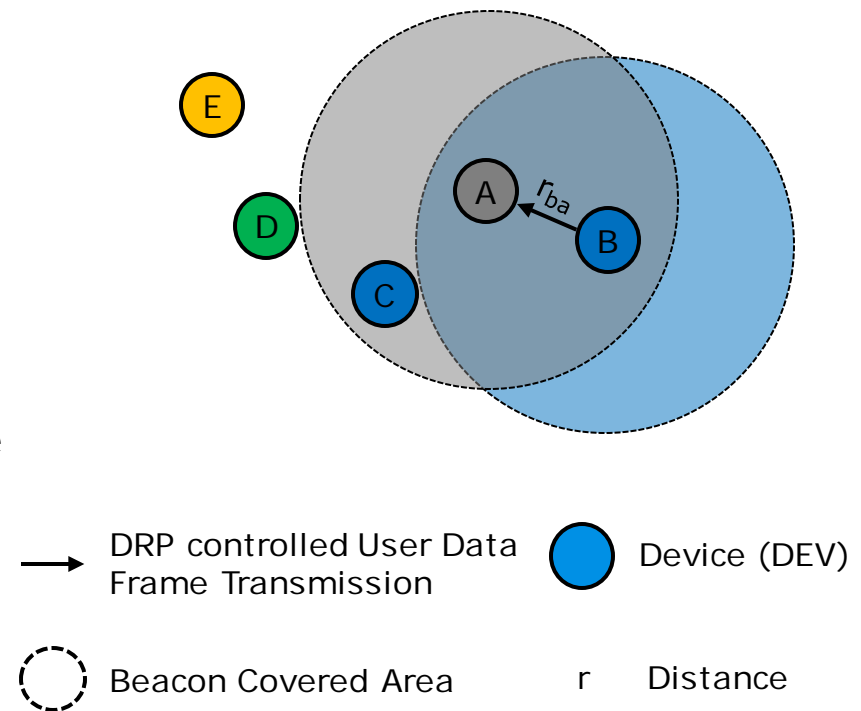
- User Data Frames are carried in reserved MASs negotiated previously
- Management information on DRP controlled medium access is carried in Beacons
- MASs are granted exclusively until DEV suspends DRPIE from Beacon transmission
- DEVs not located in one-hop BG from owner or target may reuse MASs randomly for user data frame transmission
- Maximum interference power is expected where Beacon Decoding Range covers same area [3]



[3] H. Rosier, C. Kukla. Performance Evaluation of DRPNext in ECMA-368 Wireless Personal Area Networks, " *Personal Indoor and Mobile Radio Communications (PIMRC), 2013 IEEE 24th International Symposium on*, 2013

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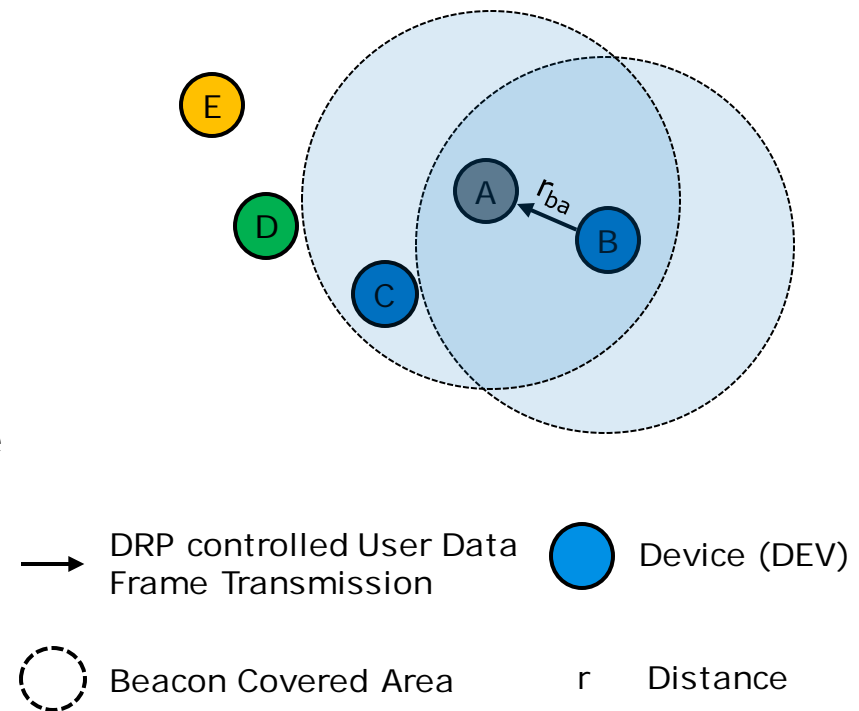
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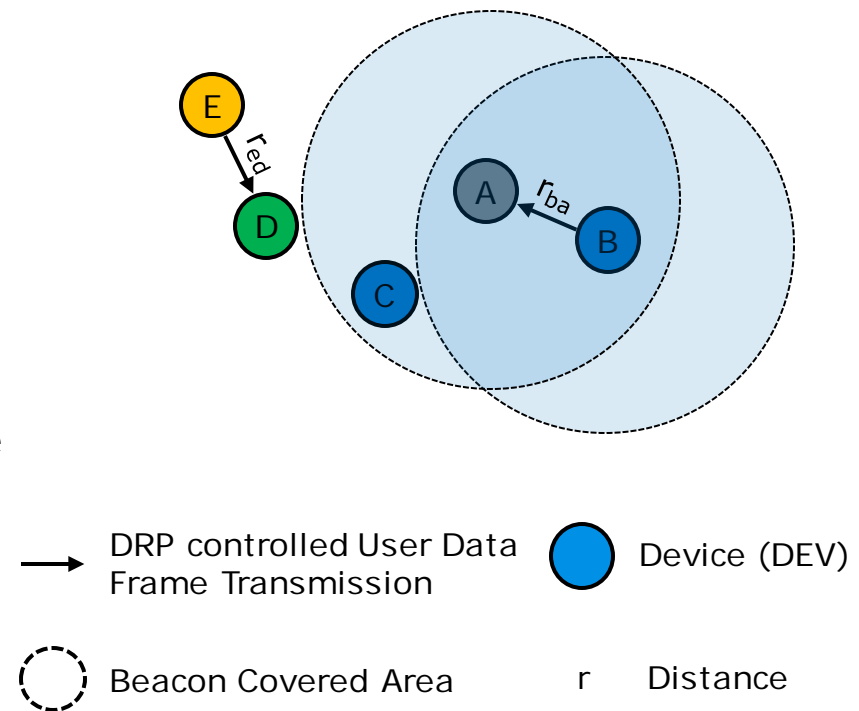
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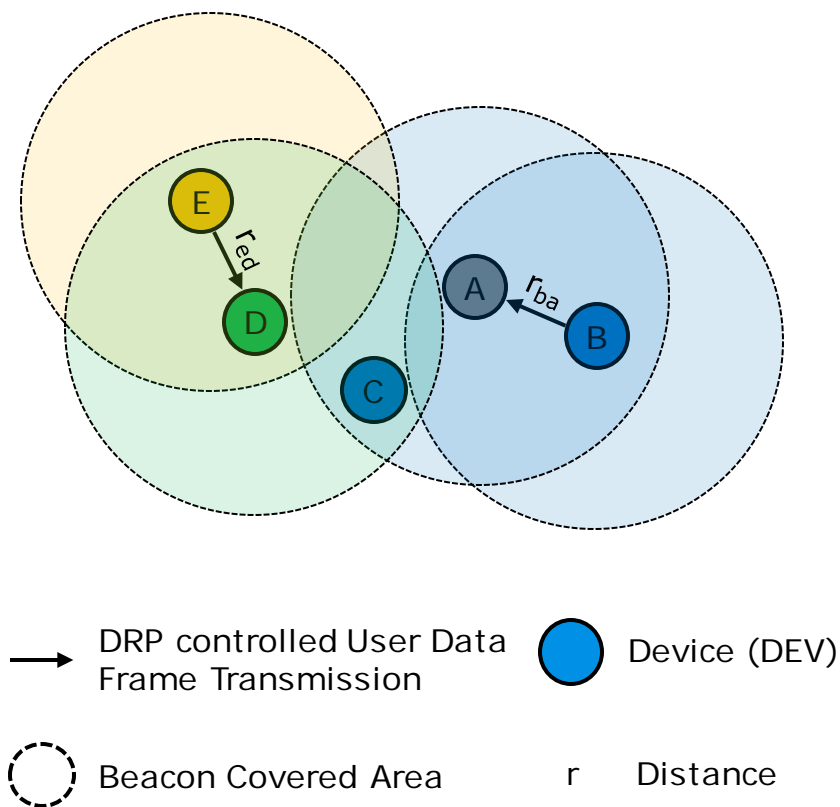
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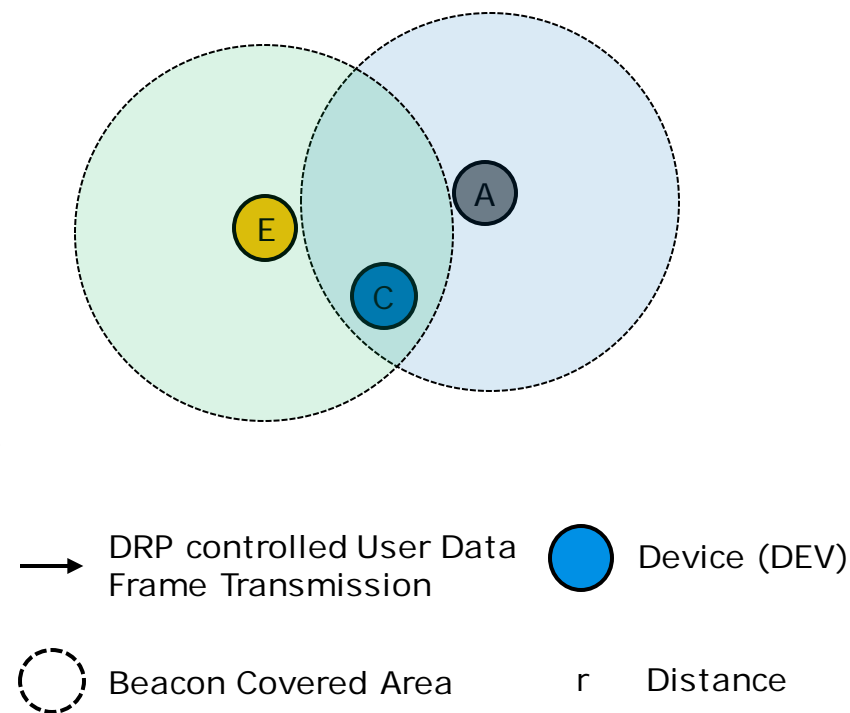
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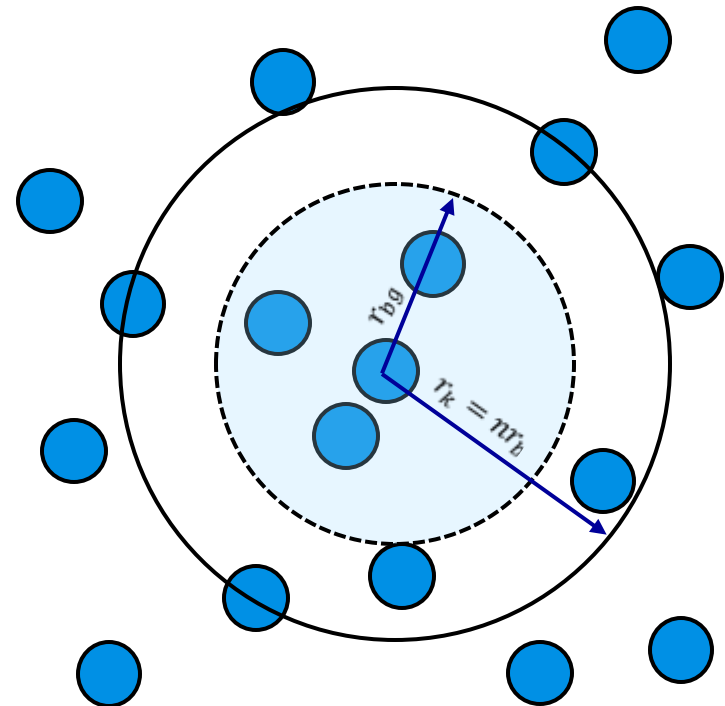
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Scenario Description & Parameters

Parameter	Value
Number of Devs per m^2	$\rho = 0.05 \frac{DEVs}{m^2}$
DEV Position	Random uniform distributed
Pathloss Calculation	Mean Pathloss Channel Model (MPCM)
Mean Beacon Decoding Range	$r_{bg} = 16m$
Frequency	3.96GHz
Data Frame Size	1500B
Modulation and Coding Schemes (MCS)	53.3MB/s ... 480MB/s
Evaluated Scenarios Size	$r_s = n \cdot r_{gb}$ $n = 1, 2, \dots, 5$
Target FER	3%
Distance between owner and target, considered for MCS probability distribution	0.1m, ..., 9m



Analytical Study Connectivity & Beacon Group

- Connectivity depends on number of DEVs located in the BG

$$c_f = \sum_{i=1}^N \frac{M_i}{N(N-1)} = \frac{\rho\pi \cdot r_{bg}^2 - 1}{\rho\pi \cdot r_s^2 - 1} \approx \frac{r_{bg}^2}{r_s^2}$$

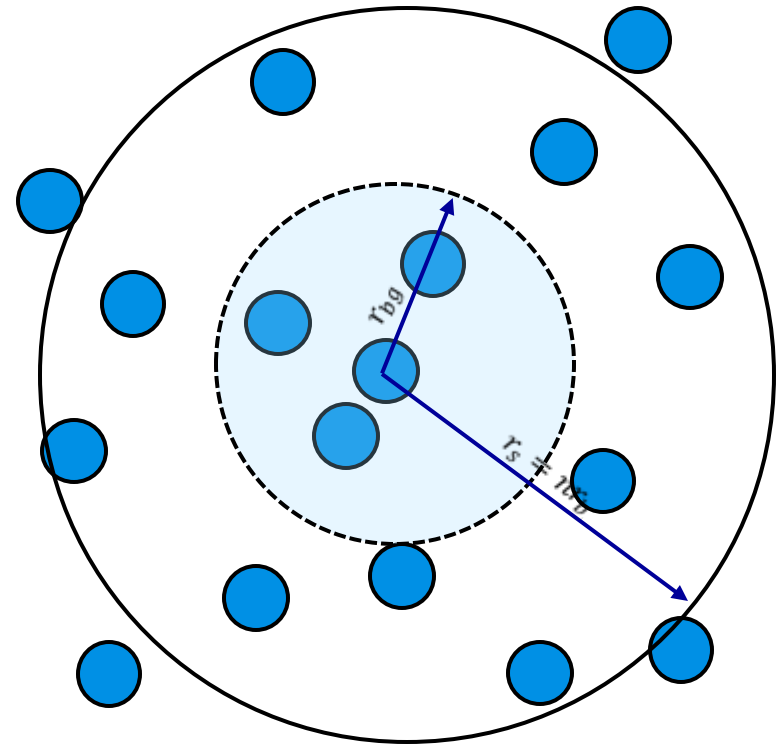
connectivity c_f , N number of DEVs located in scenario, ρ DEV density, M_i number of DEVs in one-hop BG neighborhood, r_{bg} Beacon Decoding Range, r_s evaluated scenario size

- Mean Number of DEVs located in BG

$$N_{BG} = N \cdot c_f$$

- Number of DEVs where first $k = 1$ co-channel interferer may occur

$$N_{k=1} = (1 - c_f) \cdot N$$



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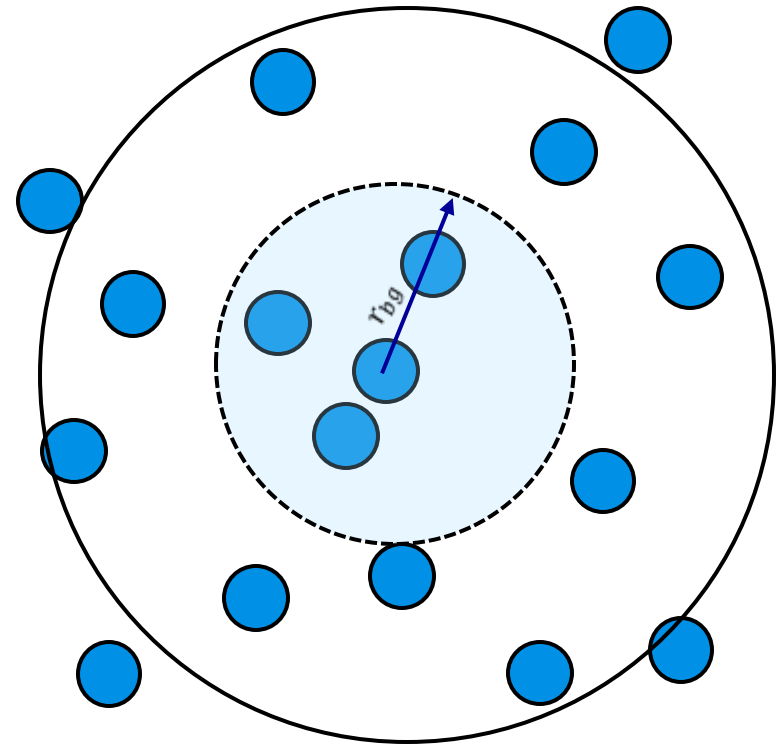
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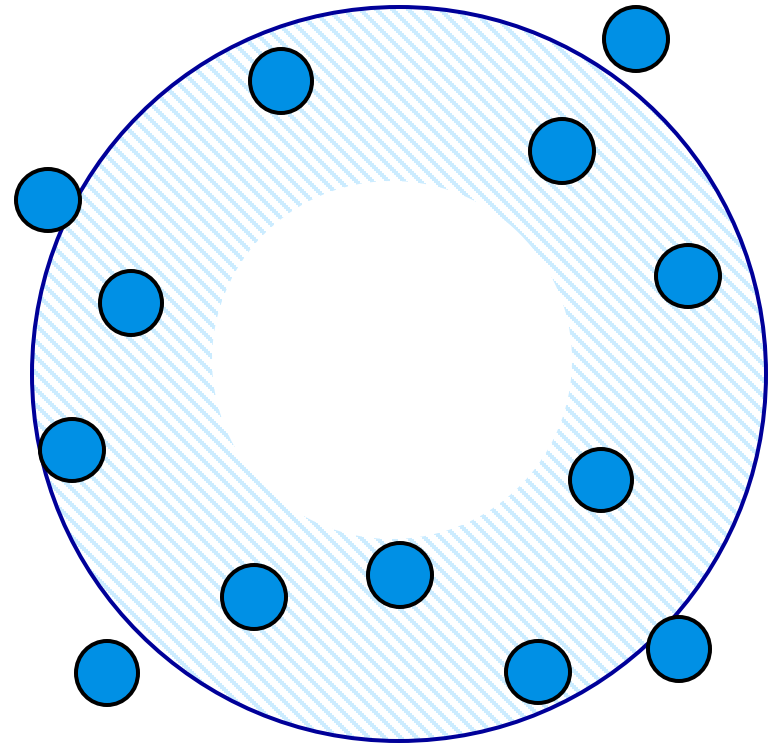
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Analytical Study

Number of DEVs able to Reuse MASs

- Number of DEVs where second $k = 2$ co-channel interferer may occur

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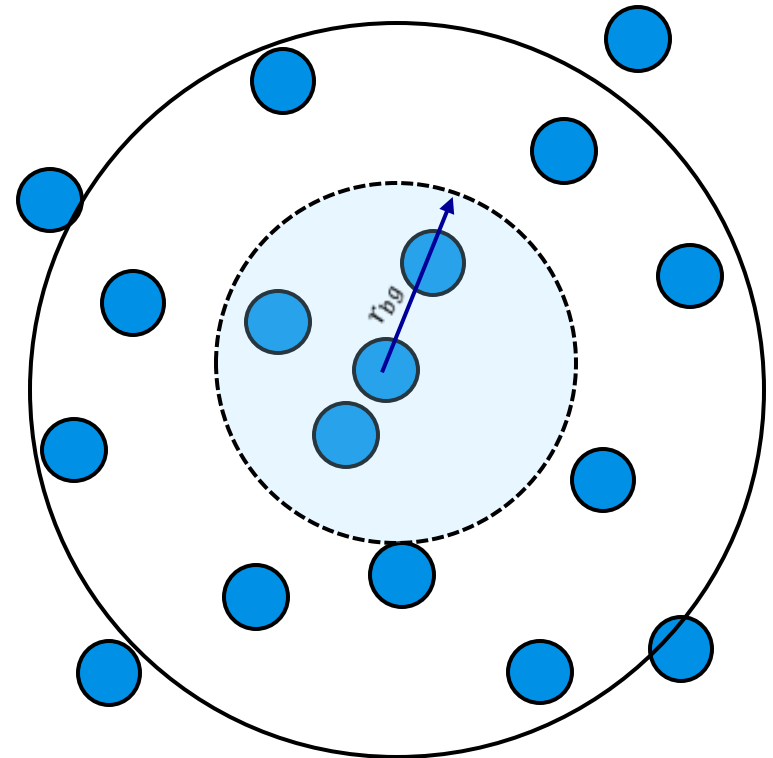
- Considers DEVs located in more than one BG

$$(N_1 - \widetilde{N}_2) \cdot c_f$$

- Leads in general to [4]

$$N_k = N(1 - c_f)^{k+1} + N_{k-1}c_f + \sum_{i=1}^{k-2} (N_i - N_{i+1}) c_f^{k-i}$$

where N_k states number of DEVs considered to calculate k^{th} reuse of MASs



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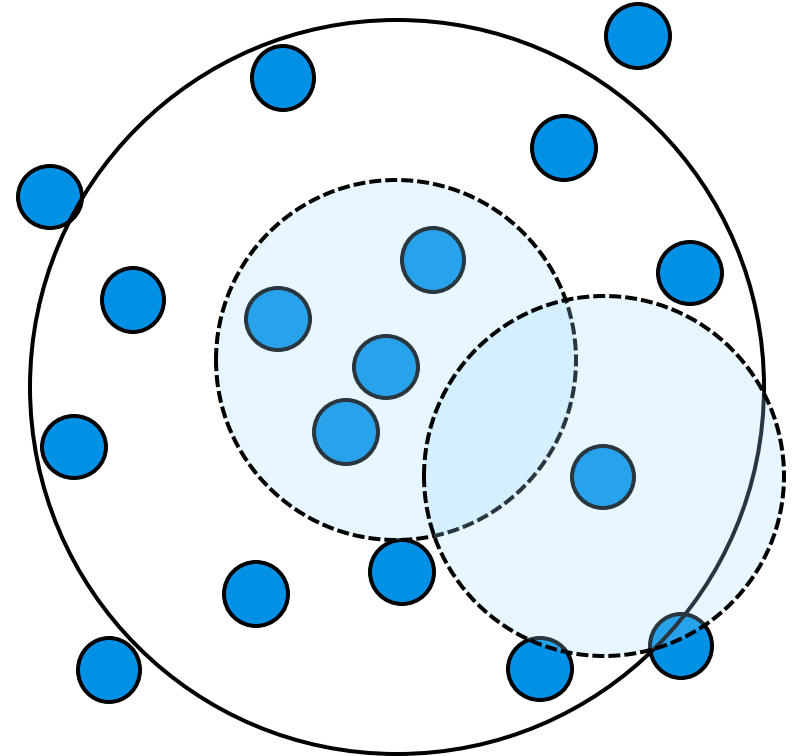
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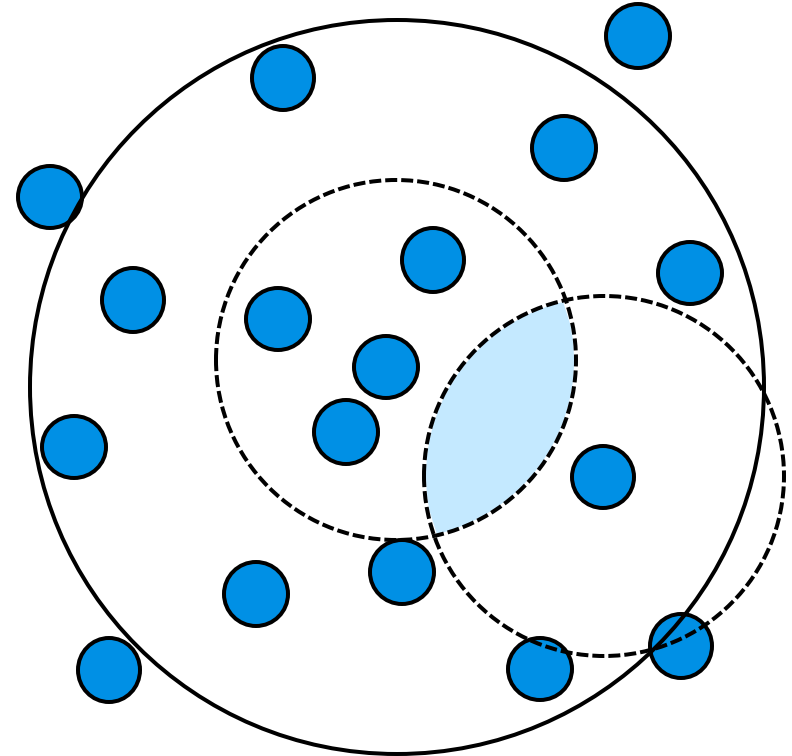
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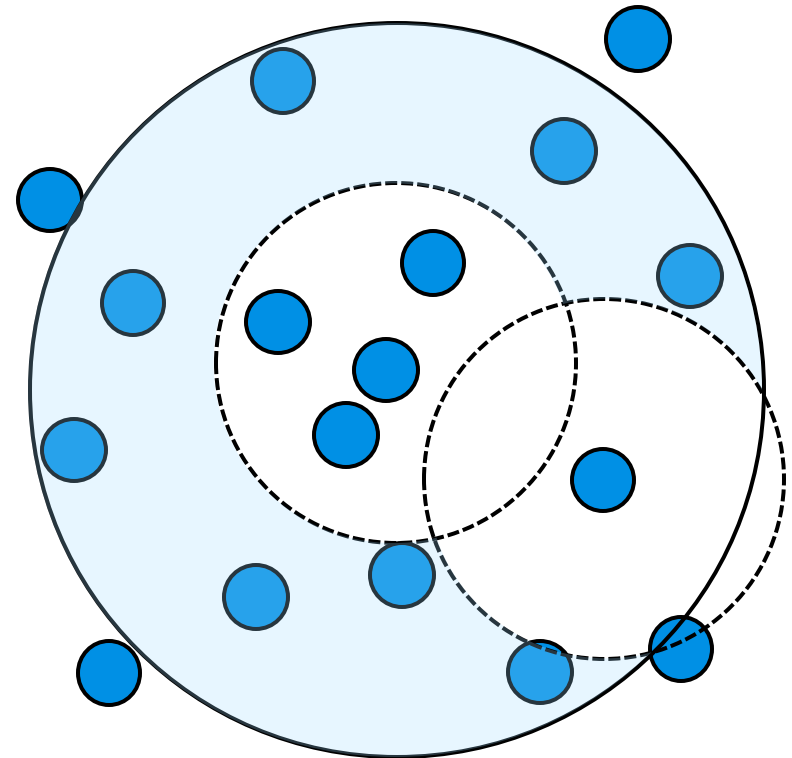
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Analytical Study

Probability of co-Channel Interferer

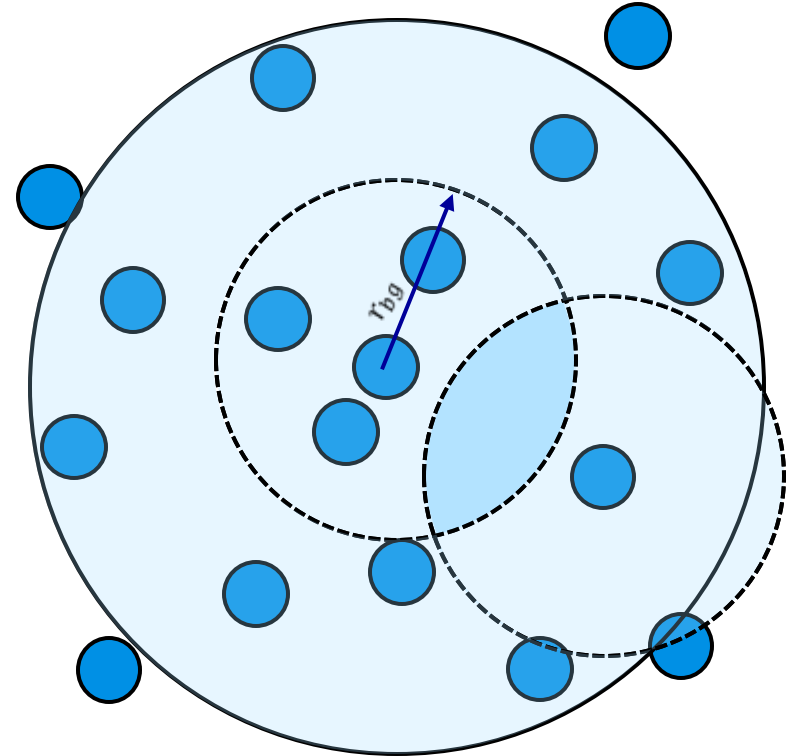
- Mean number of available MASs for DRP controlled user data frame transmission

$$\bar{d} = \frac{1}{N_{BG}} \sum_{i=0}^{N_{BG}-1} I_{MAS} - i$$

- Total number of available MASs in SF $I_{MAS} = 256$
- Probability to find exactly k co-channel interferer

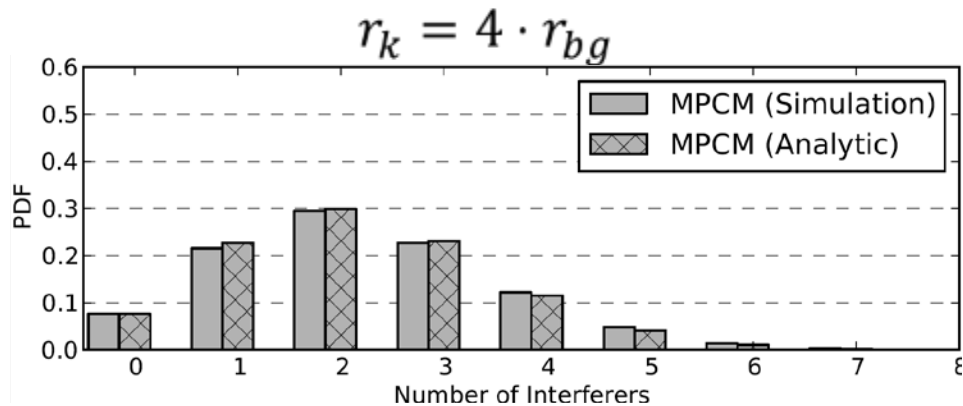
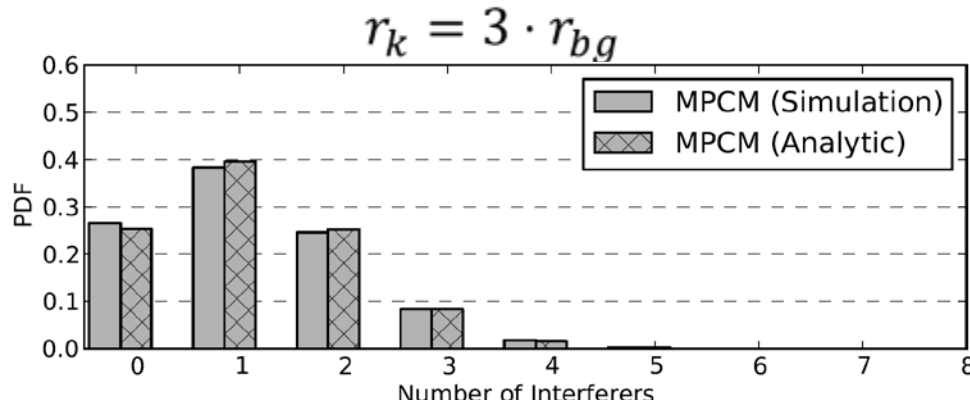
$$p_k = \binom{N_k}{k} \left(\frac{1}{\bar{d}}\right)^k \left(1 - \frac{1}{\bar{d}}\right)^{N_{k+1}}$$

$$p_0 = 1 - \sum_{k=1}^{\infty} p_k$$



Results

Probability of Co-Channel Interferer

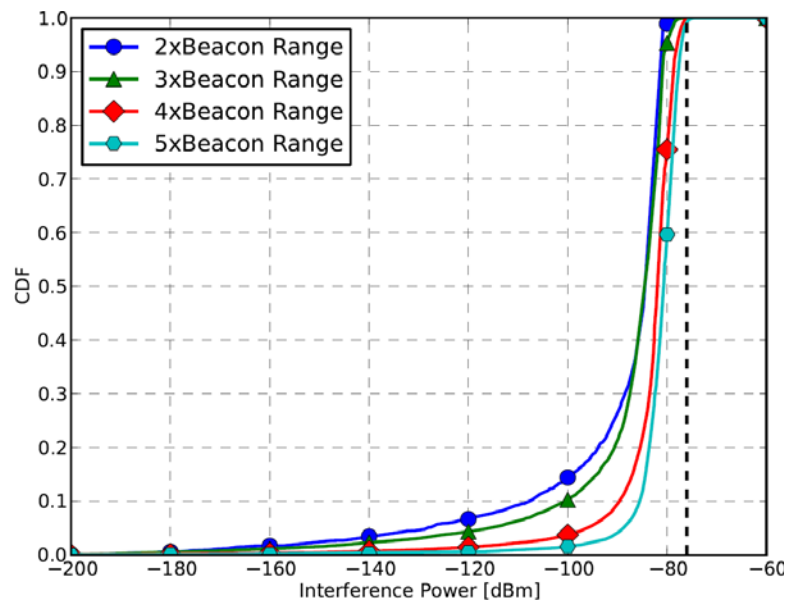


- Number of interfered DRP controlled user data frame transmissions is equivalent to MAS reuse in the scenario
- Probability that no DEV allocates MAS operated for user data frame transmission by another DEV decreases from 0.28 to 0.09 where radius r_k increases from $3 \cdot r_{bg}$ to $4 \cdot r_{bg}$
- Monte Carlo experiment complies results obtained analytically

Results

Maximum Interference Power

- Interference power for increasing scenario size regarding parameter value $r_k = n \cdot r_{bg}$ with $n = 2, \dots, 5$

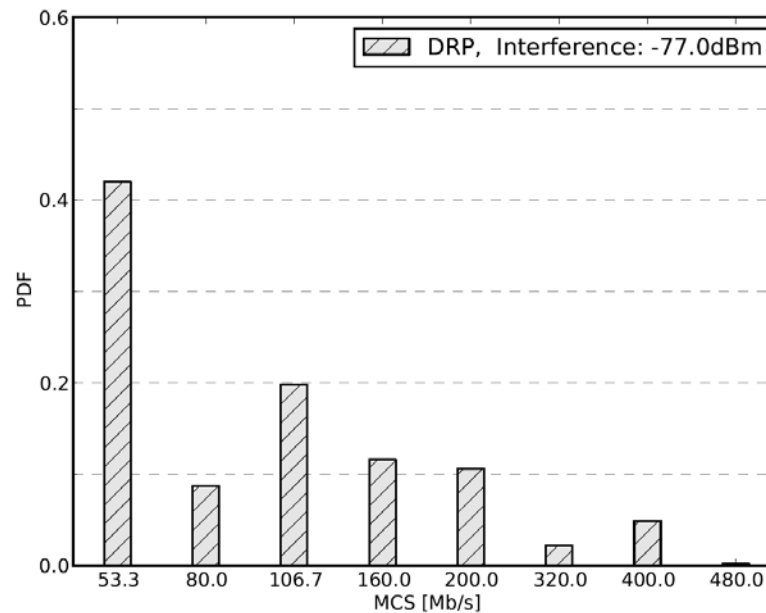


- CDF tends to step function in scenarios where r_k increases
- Concerning results DRP limits interference power to value $\approx -77dBm$

Results

MCS Probability Distribution & Achievable Throughput

- MCS probability results considers random distance between DRP reservation owner and DRP reservation from uniform distribution and constant interference power



- Achievable throughput concerning rate fair scheduling calculates from

$$t_{hr} = \left(N_{BG} \cdot \sum_{i=1}^7 \frac{p(MCS_i)}{rate_i} \right)^{-1} = 2.026 \frac{Mb}{s}$$

$p(MCS_i)$ probability to operate MCS_i , $rate_i$ data rate offered from MCS_i

- Achievable throughput in Beacon Group is obtained from $N_{BG} \cdot t_{hr} = 81.47 \frac{Mb}{s}$

Conclusion & Outlook

Conclusion:

- Interference power in UWB scenarios where DEVs operate DRP for user data frame transmission and allocate MAS randomly is limited to $-77dBm$
- Throughput achieved from DRP controlled medium access in densely populated network scenarios calculates to $81.47 \frac{Mb}{s} \frac{1}{BG}$

Outlook:

- Apply Interference Aware Scheduling to gain throughput in UWB WPAN scenarios^[5]

^[5]H. Rosier, J. Frerichs. "Interference Aware Scheduling for Ultra Wideband Networks," *Wireless and Mobile Computing, Networking and Communications (WiMob)*, IEEE 6th International Conference on, Niagra Falls 2010

Thank you for your attention!

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hor@comnets.rwth-aachen.de