

Network Planning for Stochastic Traffic Demands using a Genetic Algorithm

Dr. Nga Tran

Prof. Dr. Andreas Timm-Giel

Basic network planning problem

Given

- Network information (topologies, link cost, etc.)
- Traffic demands

Determine

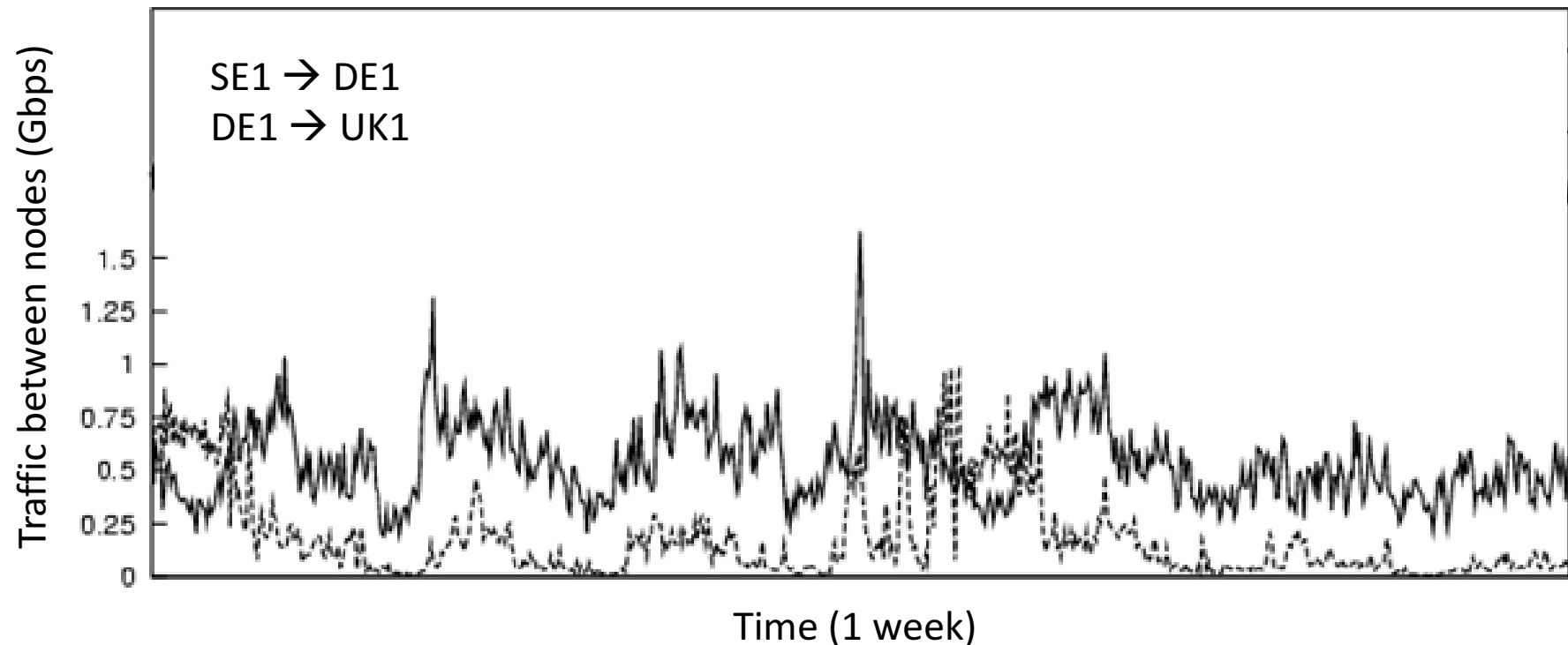
- Route for each demand
- Bandwidth allocated for each link

Objective

The total network cost
subject to network constraints is minimized.

Problem

Traffic demands are not deterministic but stochastic.



How to handle the uncertainty?

Planning Approaches

- Mean-rate based
- Pro: low cost
- Con: high amount of traffic is discarded
- Peak-rate based
- Pro: able to handle a large traffic demand variation
- Con: very expensive

Statistical network planning

- Objective: minimize the cost
- Constraint: accept a violation probability for the link load (Grade of Service – GoS)

Mathematical model (1)

- Parameters:
 - Traffic demand distribution of a nodepair (sd) : t^{sd}
(Traffic demands are statistically independent.)
 - Cost of a bandwidth unit on the link ij : c_{ij}
 - GoS agreement: ε
- Variables:
 - Routing (binary) decision variable: f_{ij}^{sd}
 - Bandwidth assignment variable: b_{ij}

Mathematical model (2)

- Objective: minimize the total cost

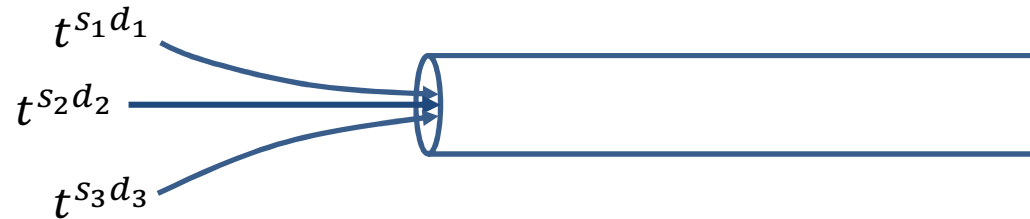
$$\min \sum_{ij} c_{ij} \cdot b_{ij}$$

- Multi-commodity flow constraint

$$\sum_j f_{ij}^{sd} - \sum_j f_{ji}^{sd} = \begin{cases} 1 & i = s \\ -1 & i = d \\ 0 & i \neq s, d \end{cases} \quad \forall i$$

Mathematical model (3)

- Capacity constraint



Aggregated traffic distribution:

$$t = t^{s_1 d_1} \otimes t^{s_2 d_2} \otimes t^{s_3 d_3}$$

Probability that the aggregated traffic is lower than the capacity b of a link is: $\int_0^b t(x) dx$

Mathematical model (4)

- Capacity constraint

$$\int_0^{b_{ij}} \left(f_{ij}^{s_1 d_1} \cdot t^{s_1 d_1} \otimes f_{ij}^{s_2 d_2} \cdot t^{s_2 d_2} \otimes \dots \otimes f_{ij}^{s_n d_n} \cdot t^{s_n d_n} \right) (x) dx \geq 1 - \varepsilon \quad \forall i, j$$

$$\longrightarrow \prod_{\{s, d; f_{ij}^{sd} = 1\}} \int_0^{b_{ij}} f_{ij}^{sd} \cdot t^{sd} (x) dx \geq 1 - \varepsilon$$

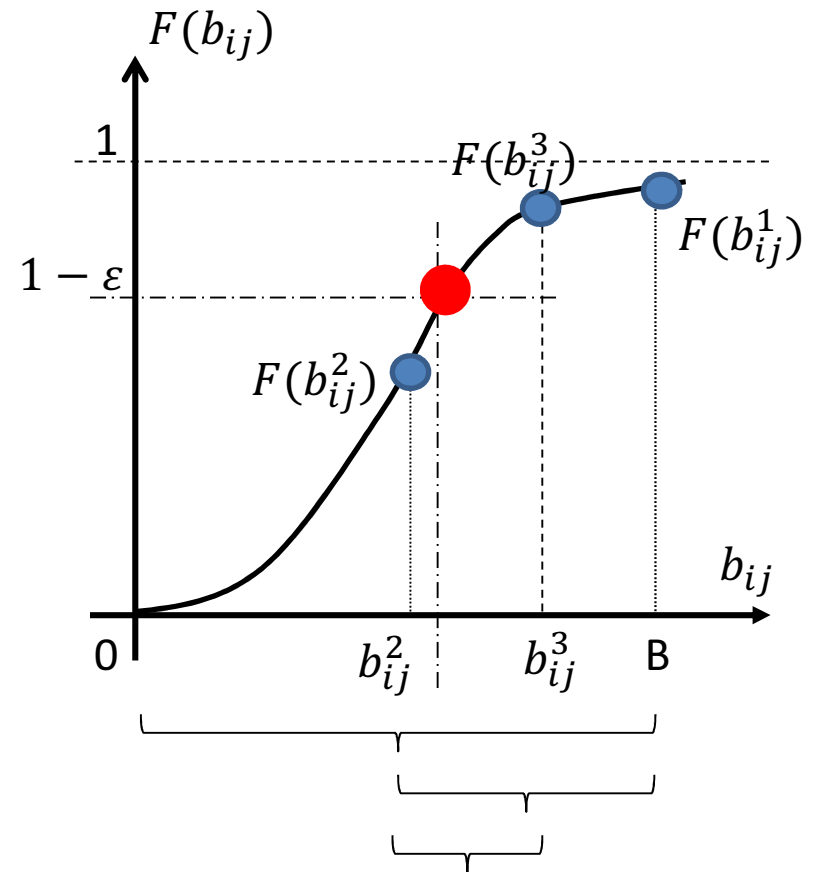
$$P\{\text{link } (i, j) \text{ is overloaded}\} \leq \varepsilon$$

$$b_{ij} = ?$$

Calculate the link bandwidth b_{ij}

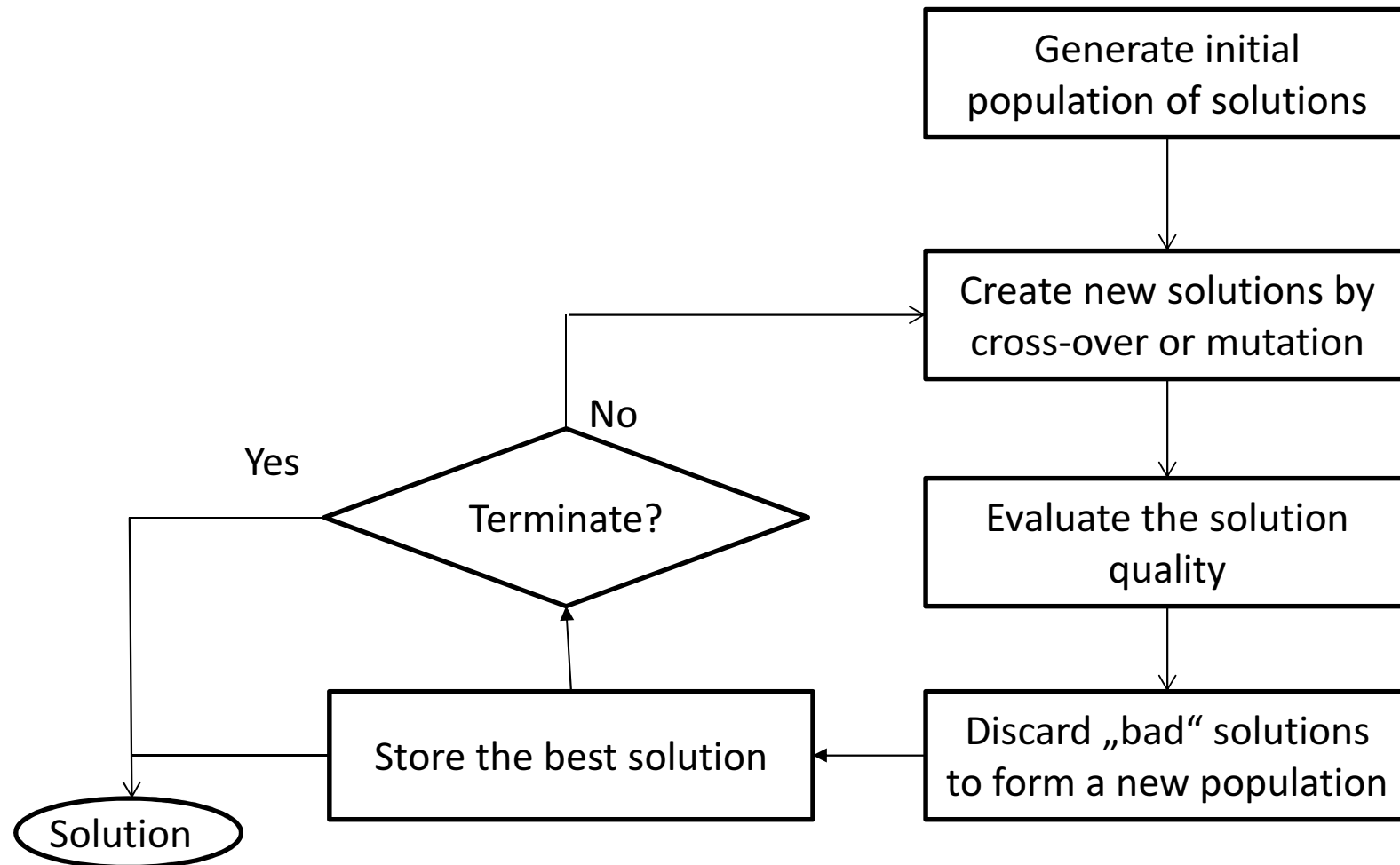
$$F(b_{ij}) = \prod_{\{s,d;f_{ij}^{sd}=1\}} \int_0^{b_{ij}} f_{ij}^{sd} \cdot t^{sd}(x) dx \geq 1 - \varepsilon$$

- $F(b_{ij})$ is the CDF of the traffic distribution on the link
→ a monotone function
- bisection method can be used to find b_{ij}



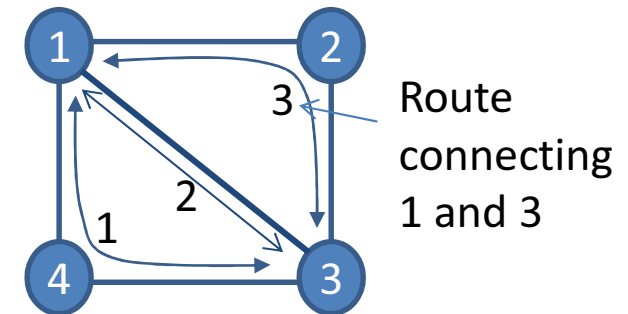
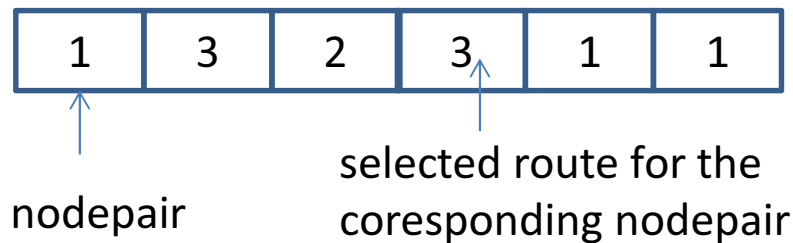
Solution interval

Genetic algorithm (1)



Genetic algorithm (2)

- A set of routes for each nodepair is pre-selected.
- A solution is encoded in a chromosome

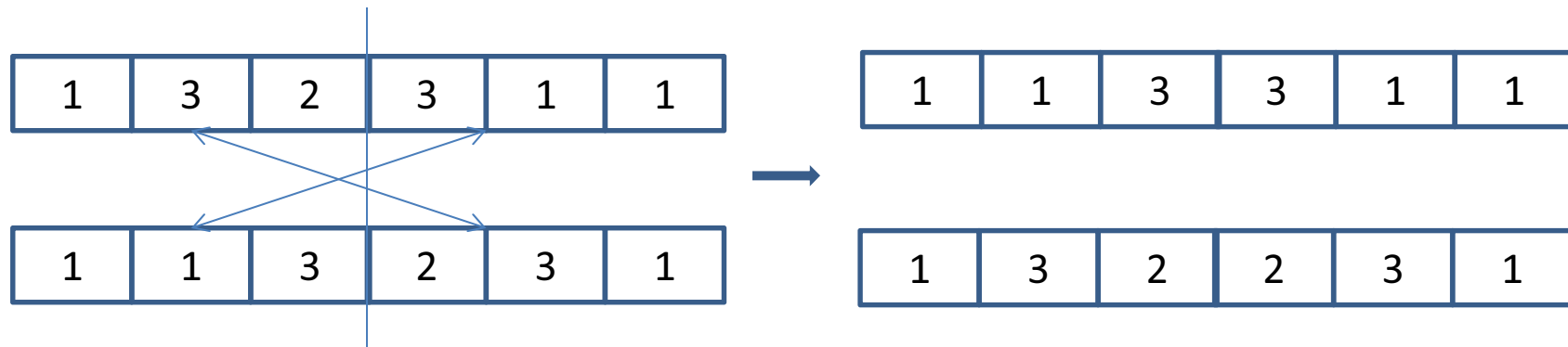


→ a chromosome represents a routing solution for the problem.

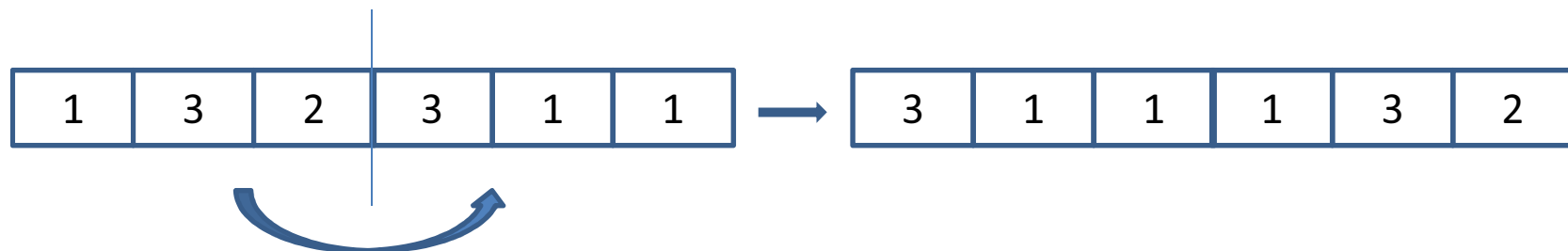
- Calculate the required bandwidth for each „chromosome“ to guarantee GoS
→ calculate the cost of each „chromosome“

Genetic algorithm (3)

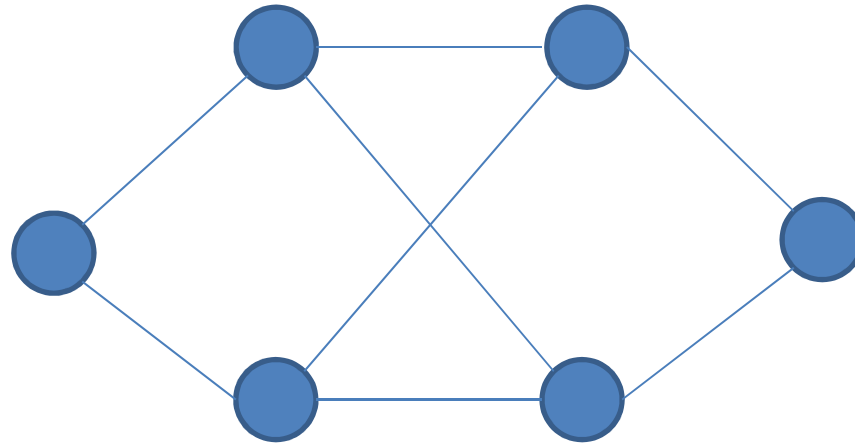
- Cross-over



- Mutation



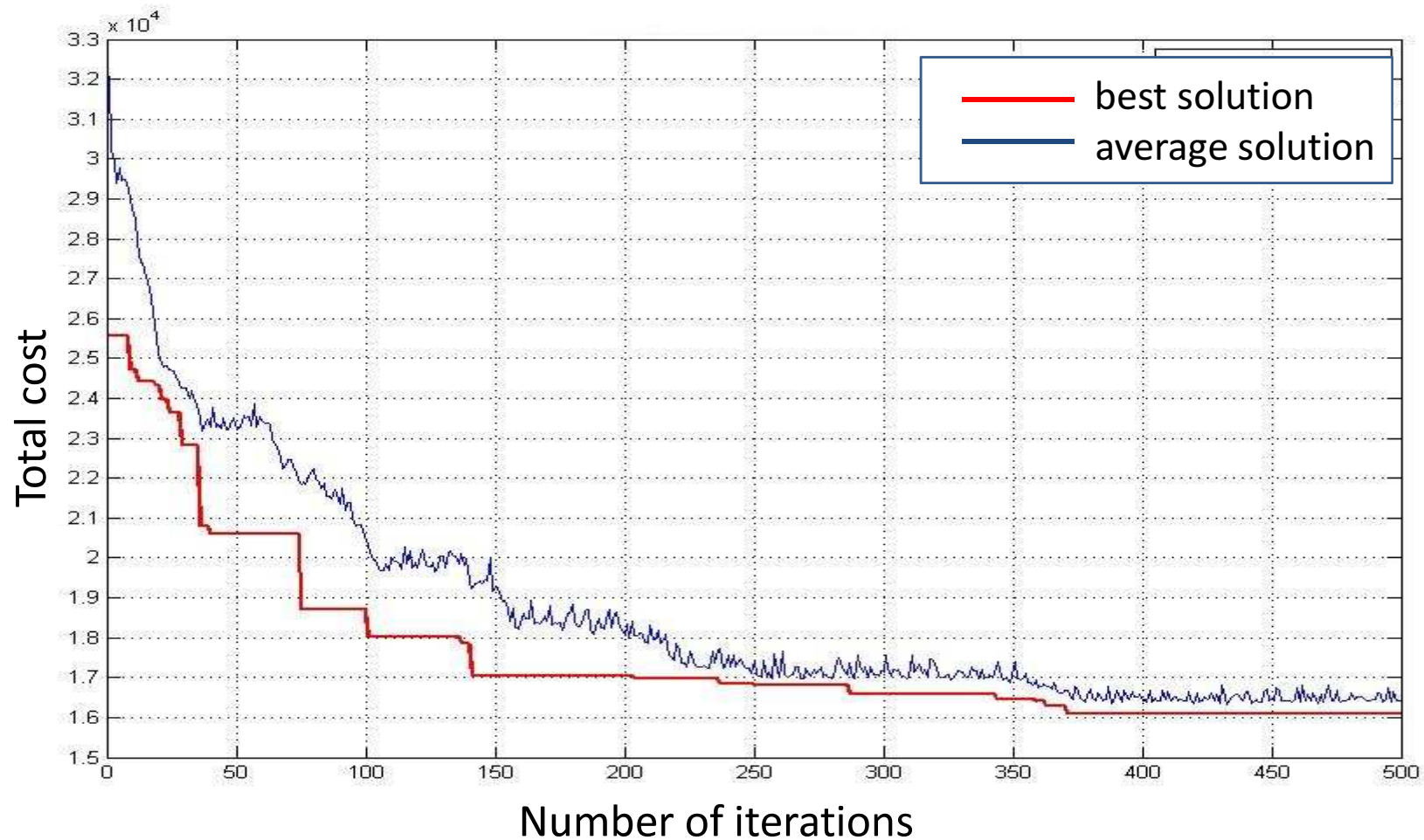
Performance Evaluation



- Traffic demands are taken from GEANT project*
- Link cost = 1 per Mbps
- GoS agreement: $\varepsilon = 5\%$

(*) <http://www.geant.net/>

Performance of the genetic algorithm

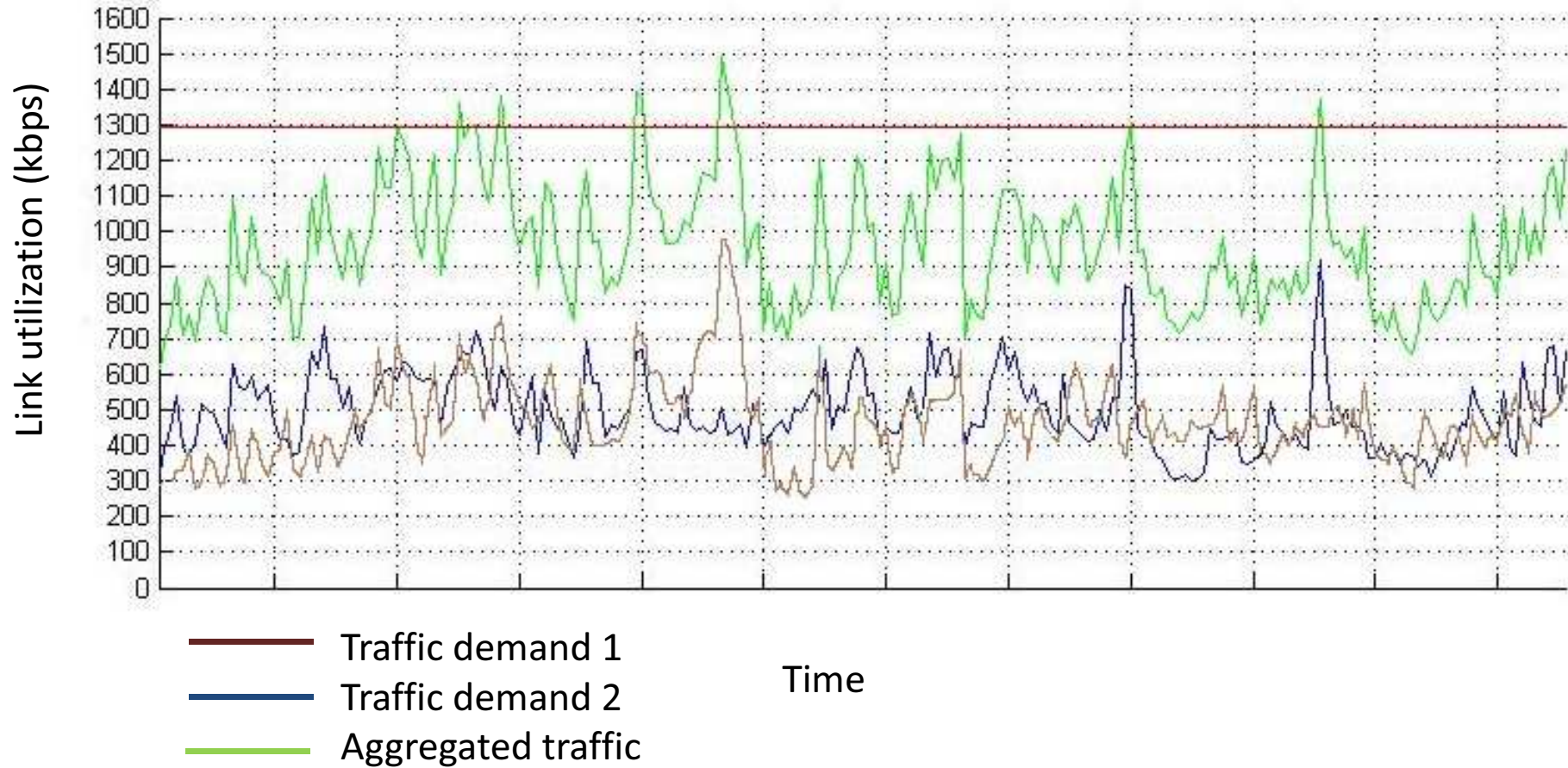


Results

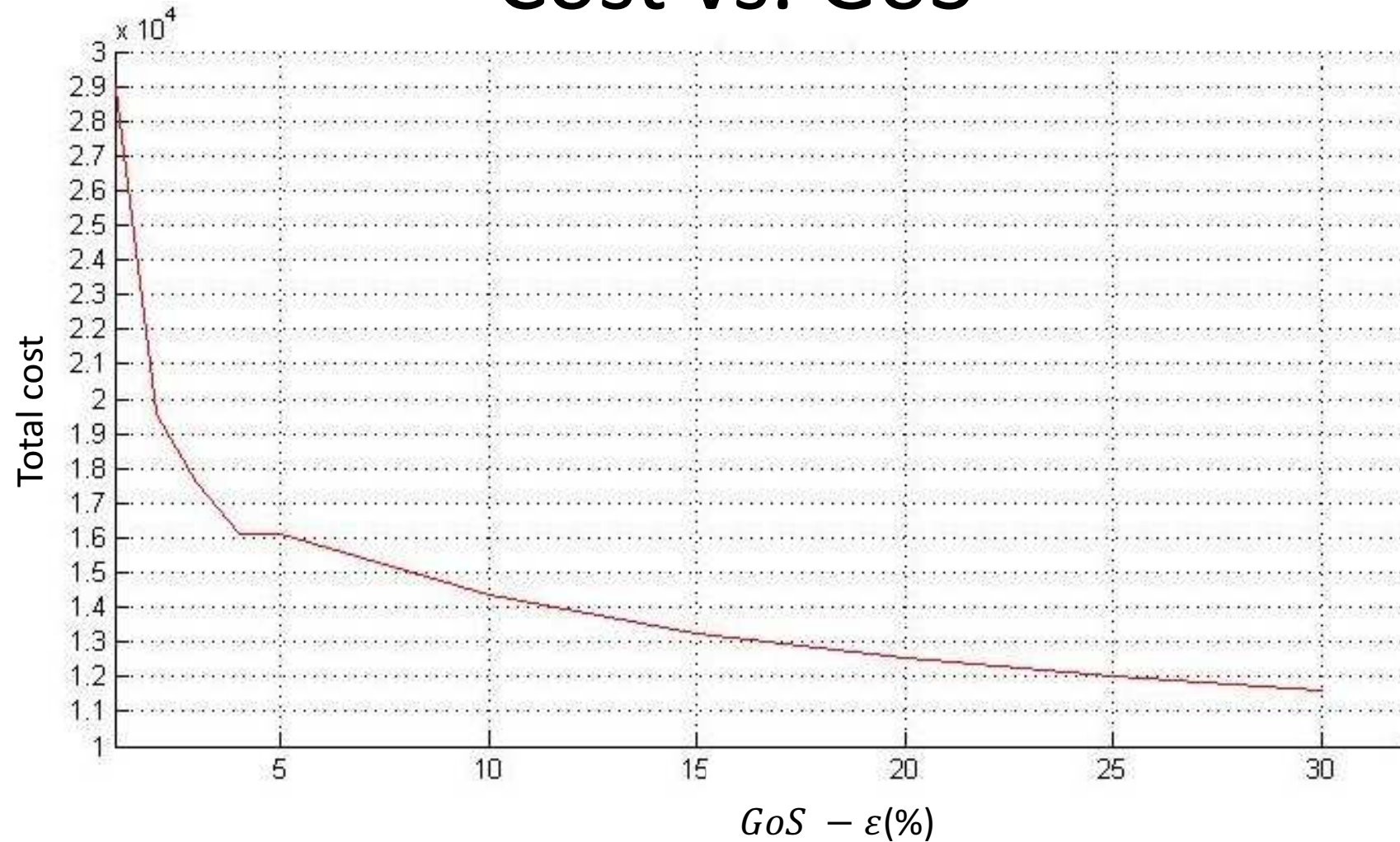
	Mean rate	Peak rate	Stastical
Normalized cost	0.42	1	0.53
Link overload probability	18% - 25%	0%	5% (ϵ)

- Peak rate planning: too expensive
 - Mean rate planning: high link overload probability
- Statistical planning is a good compromise

Link utilization



Cost vs. GoS



Conclusion

- Proposal of a solution for network planning problem with stochastic traffic demands using a genetic algorithm.
- Limitation: guarantee the overload probability for each link only
- Future work: guarantee the GoS for each end-to-end flow