

Switching of Routing Algorithms in Wireless Networks for Fire Fighting

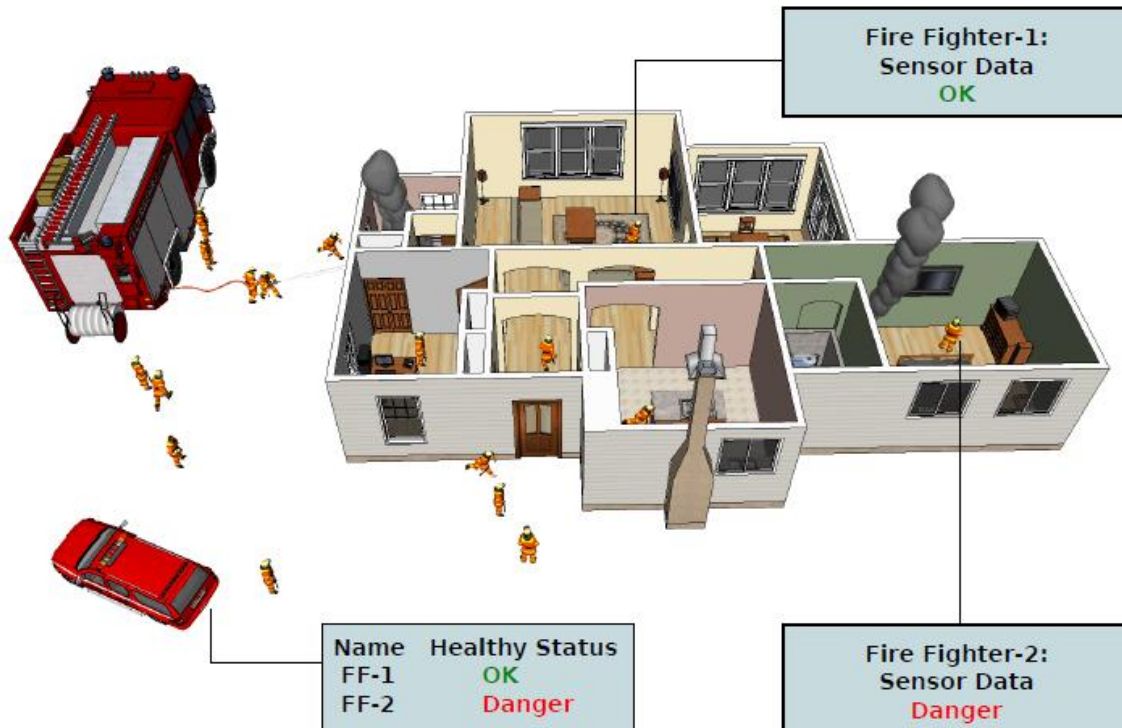
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Oct. 19, 2012

Outline

- Problem Statement
- Proposed Adaptive Routing Framework
- Analytical Model for Switching of Routing Algorithms
- Conclusions and Future Work

Fire Fighting Scenario

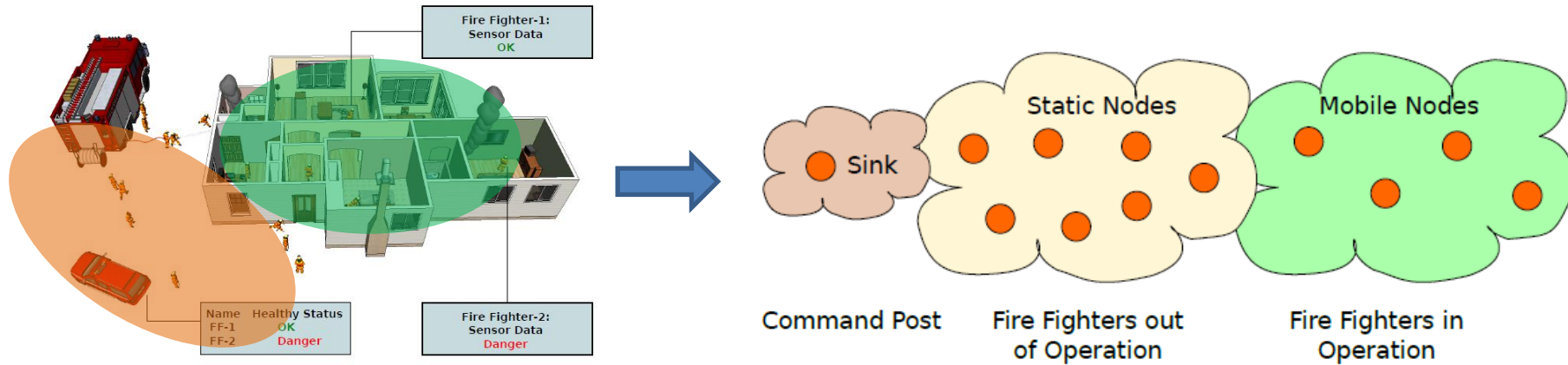


- A building is on fire
- Some fire fighters are in operation
- Rest of the fire fighters stay outside as backup
- All sensor data should be sent to the command post

Classification of Routing Algorithms

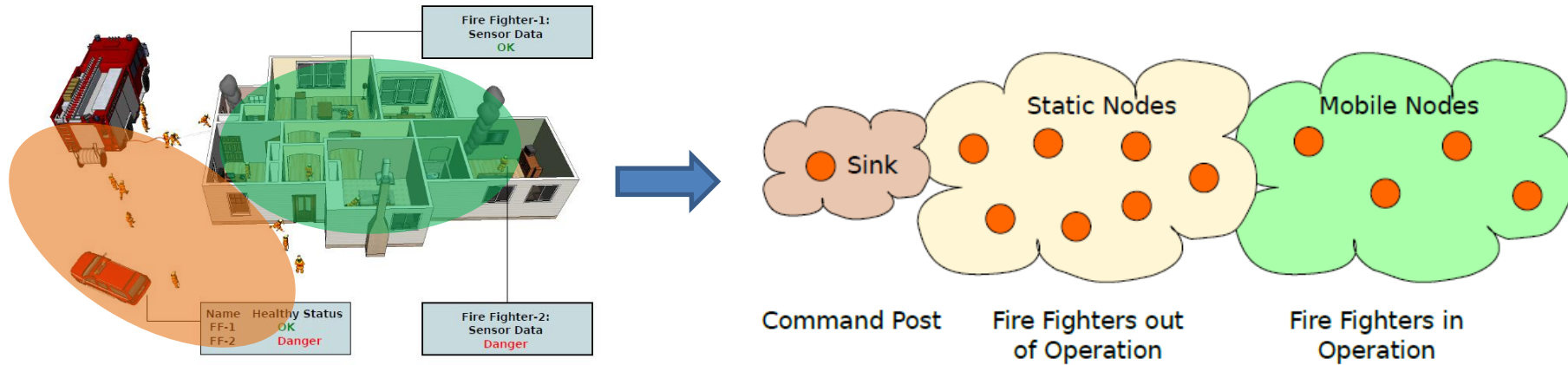
- Mobile Ad-hoc Network (MANET)
- Proactive
 - Maintain routes all the time
 - Difficult to cope with mobility
- Reactive
 - On demand
 - Inefficient if the traffic is low and route response time is important
- Hybrid
 - Try to integrate the advantages of both proactive and reactive routing

Problem Statement



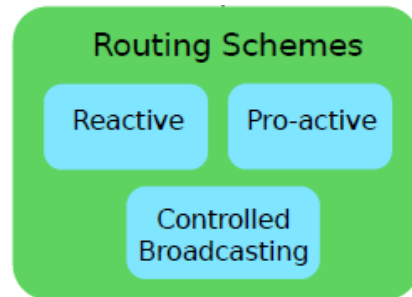
- Heterogeneous network
- Objective:
 - Best routing performance
- Static nodes: low dynamic \rightarrow stable connectivity \rightarrow proactive routing, like [Optimized Link State Routing \(OLSR\)](#)
- Mobile nodes: high dynamic \rightarrow instable connectivity \rightarrow reactive routing, like [Ad-hoc On-demand Distance Vector \(AODV\)](#)
- Solution:
 - A [self adaptive](#) routing protocol, which allows [coexistence](#) of different routing algorithms in a network.

Focus of This Presentation

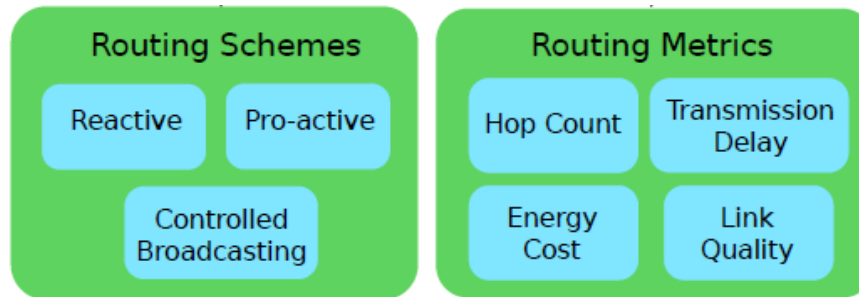


- Propose a self adaptive routing framework
- Develop decision algorithms based on analytic performance evaluation models of ad hoc routing protocols from literature

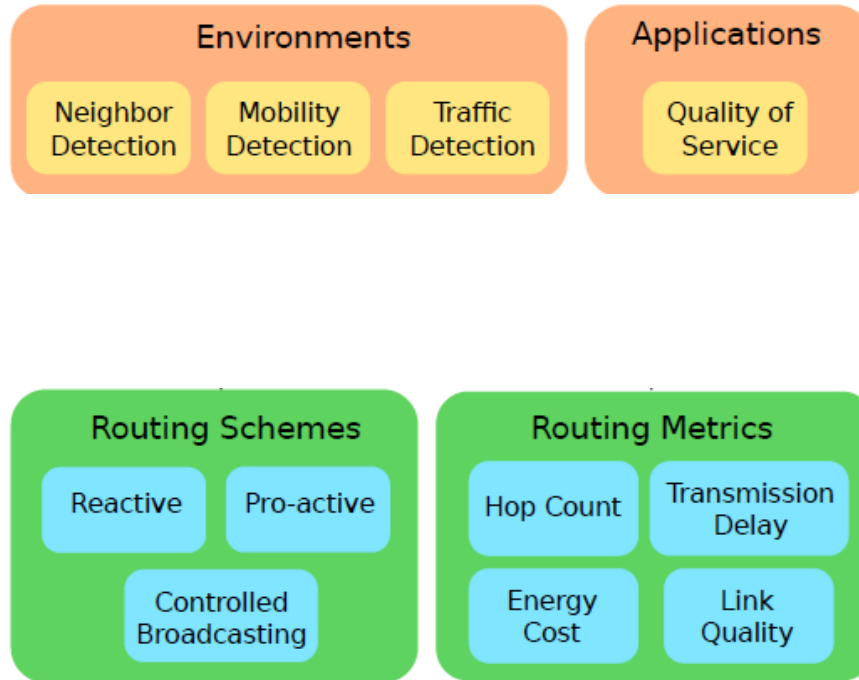
Proposed Adaptive Routing Framework



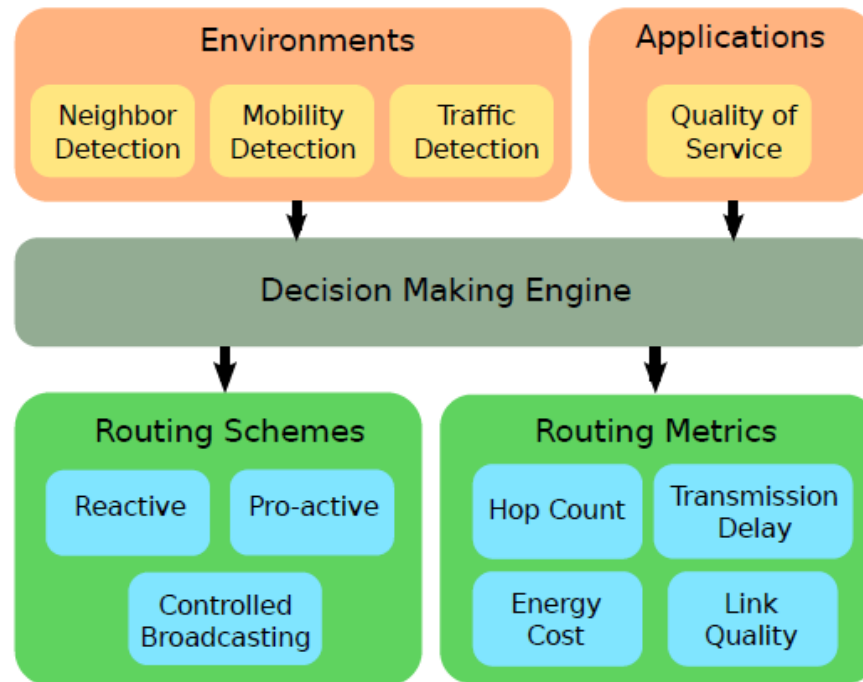
Proposed Adaptive Routing Framework



Proposed Adaptive Routing Framework



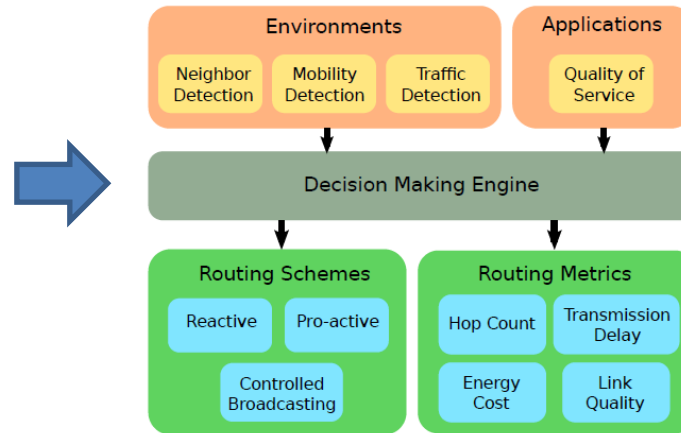
Proposed Adaptive Routing Framework



- Reactive routing as the default algorithm
- Each node may apply different routing algorithms to different neighbors
 - Through monitoring the change of its neighbor list
 - Frequent change means high dynamic → reactive routing
 - Infrequent change implies low dynamic → proactive routing

Challenge: When to switch?

Cost Functions for Switching of Algorithms



Objective: min(Cost)

$$Cost = \alpha Cost_{switch} + \beta Cost_{Algorithm}$$

$$Cost_{switch} = f_s^{Switch}(Overhead_{Switch}^{CtrlMsg})$$

$$Cost_{Algorithm} = f_s^{Algorithm}(Overhead_{Algorithm}^{CtrlMsg}(t))$$

(1) Reduce the overall cost

(2) Cost for algorithm and cost for switching

(3) Switching cost in terms of number of control messages

(4) Routing algorithm cost in terms of number of control messages

Future: other metrics are to be considered, such as packet loss rate, end to end delay, etc.

CostAlgorithm: Control Traffic Overhead (1/2)

	Reactive	Proactive
Fixed	$\lambda o_r N^2$	$h_p N + o_p t_p N^2$
Mobility	$o_r \mu a L N^2$	$o_p \mu A N_p N^2$

Network Parameters	
N	Number of nodes
μ	Link break rate (mobility)
L	Average length of a route

Data Traffic Parameters	
λ	Route creation rate per node
a	Average number of active routes per node (activity)

Reactive Protocol Parameters	
h_r	Hello rate (0 when possible)
o_r	Route request optimization factor

Proactive Protocol Parameters	
h_p	Hello rate
t_p	Topology broadcasting rate
o_p	Broadcast optimization factor
$A N_p$	Average number of active next hops

[1] Viennot, L., Jacquet, P., Clausen, T.H.: "Analyzing Control Traffic Overhead versus Mobility and Data Traffic Activity in Mobile Ad-hoc Network Protocols", *Wirel. Netw.* 10(4), 447455 (Jul 2004)

CostAlgorithm: Control Traffic Overhead (1/2)

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Mobility	$o_r \mu a L N^2$	$o_p \mu A N_p N^2$

$$O_r \lambda N N$$

Number of route requests

Every node broadcasts it once

Reduce number of broadcasting

Reactive Protocol Parameters

h_r Hello rate (0 when possible)

o_r Route request optimization factor

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Mobility	$\rightarrow o_r \mu a L N^2$	$o_p \mu A N_p N^2$

$$o_r \mu a L N$$

Number of routes
 Number of active links
 Number of link breaks
 Each node broadcasts each link break once

Reactive Protocol Parameters

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Mobility	$o_r \mu a L N^2$	$o_p \mu A N_p N^2$

$$h_p N + o_p t_p N N$$

Each node sends hello messages
 Each node broadcasts its topology
 Topology is broadcasted through the network

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$$o_p \mu A N_p N N$$

Number of active links

Number of link breaks

Each node broadcasts the link break once

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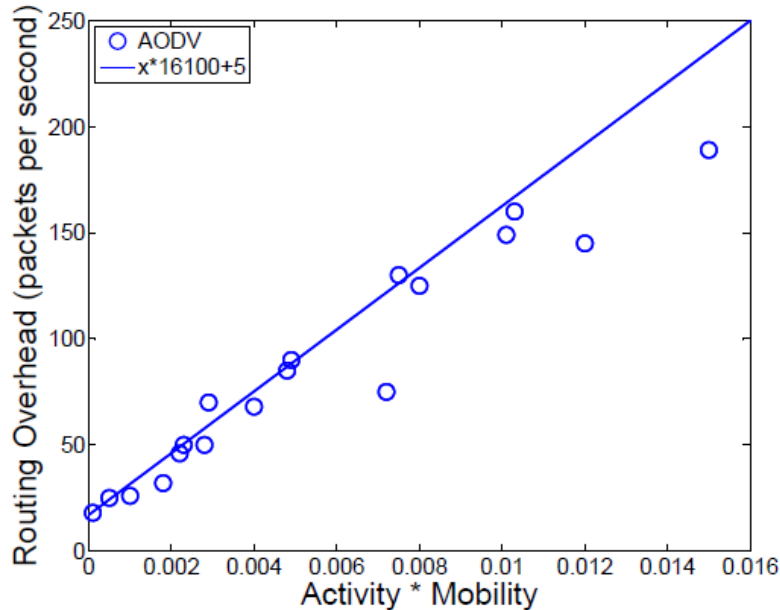
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CostAlgorithm: Control Traffic Overhead (2/2)

- Scenario
 - 50 Nodes in the area of 1500x300 meter
 - Random Waypoint mobility, different pause time
→ Mobility
 - Different number of source nodes (10, 20 or 30)
→ Activity
- Assumptions
 - No congestion occurs in the network
 - Average link break rate is constant
 - The network always remains connected

[2] J. Broch, D.A. Maltz, D.B. Johnson, Y.-C. Hu, and J. Jetcheva. A performance comparison of multi-hop wireless ad hoc network routing protocols. In MobiCom'98, October 1998. Dallas.

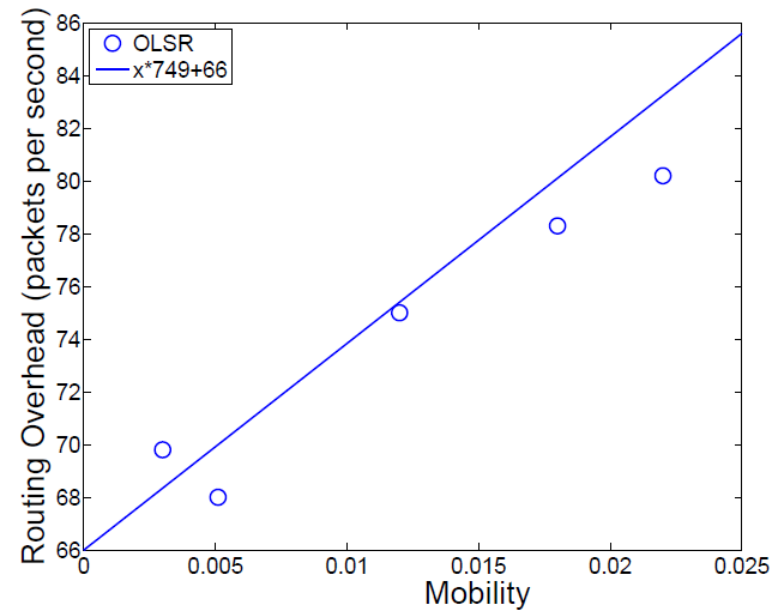
Parameter Determination



AODV: Control packets per second vs. μa

AODV

- $o_r \mu a L N^2 + \lambda o_r N^2 = 16100x + 5$
- $x = \mu a$
- $o_r = 16100 / (L N^2)$



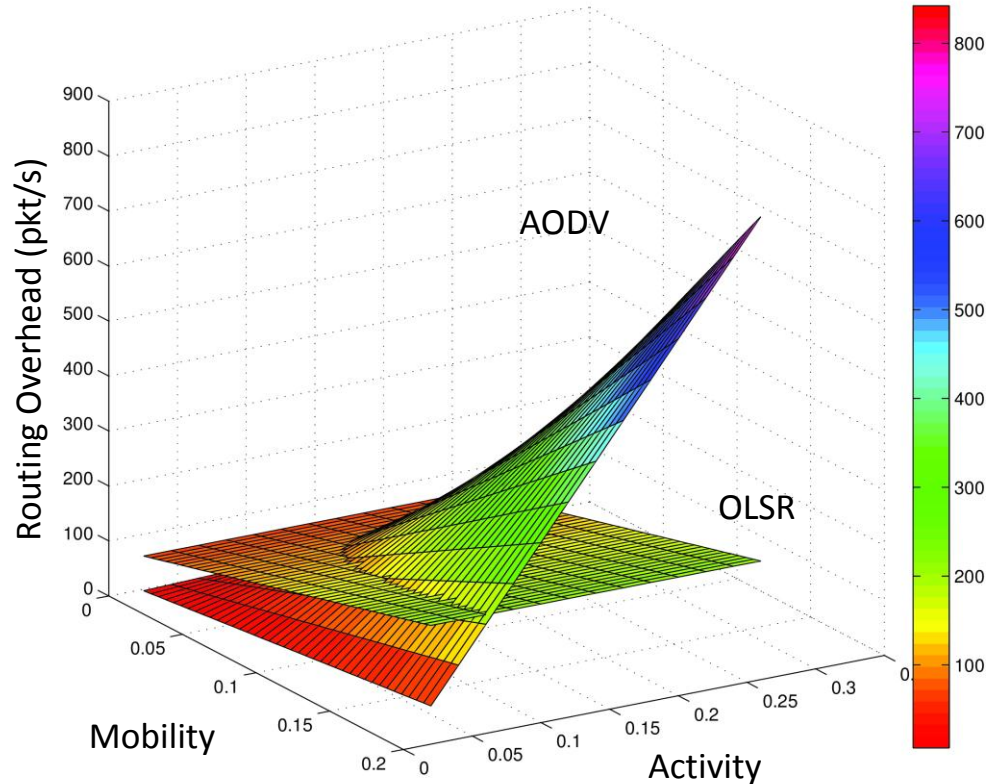
OLSR: Control packets per second vs. μ

OLSR

- $o_p \mu A N_p N^2 + h_p N + o_p t_p N^2 = 749 + 66$
- $x = \mu$
- $o_p A N_p = 749 / N^2$
- $o_p = 66 / (t_p N^2)$

Routing Overhead Comparison - Analytic

- Routing overhead vs. mobility and activity
- AODV has lower overhead when activity is low
- OLSR outperforms AODV when activity increases



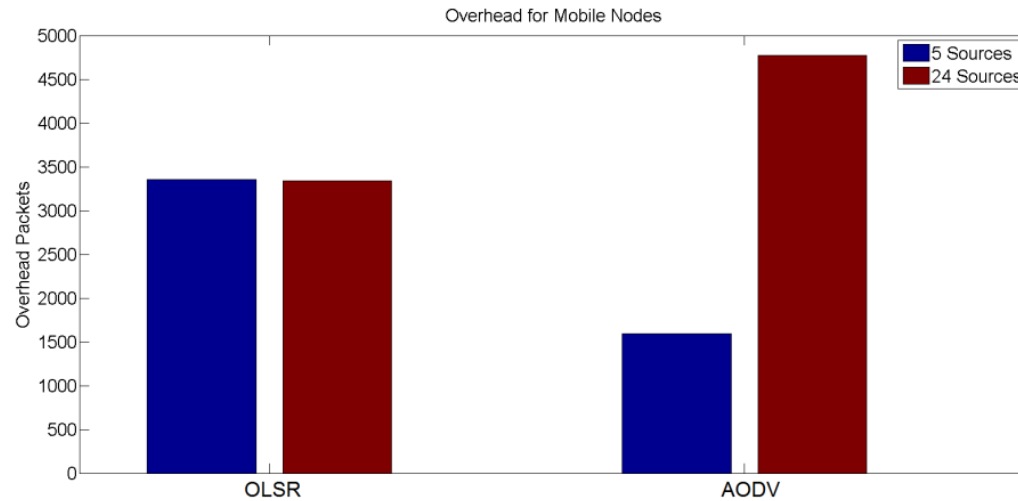
AODV

- $o_r \mu \alpha L N^2 + \lambda o_r N^2$

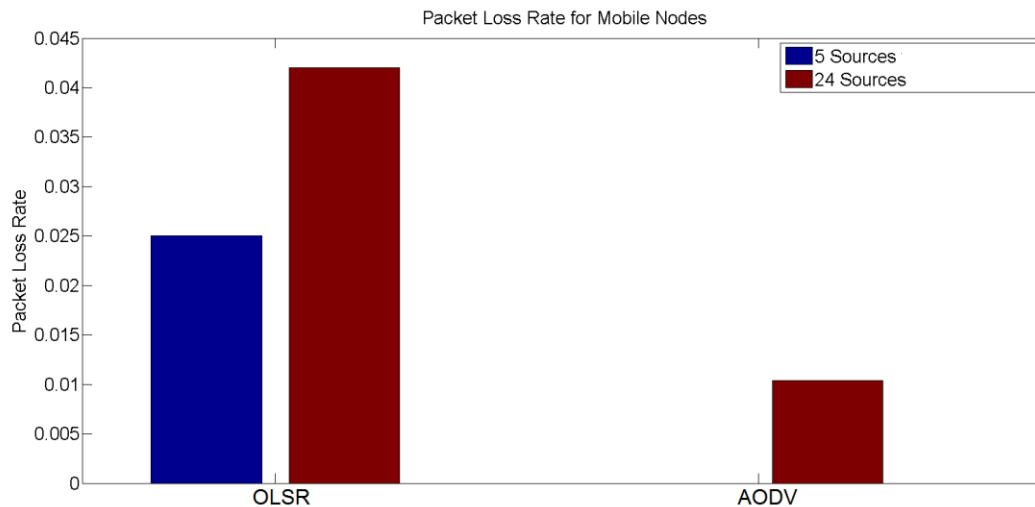
OLSR

- $o_p \mu A N_p N^2 + h_p N + o_p t_p N^2$

Routing Overhead Comparison - Simulative



Control Overhead with Mobility



Packet Loss Rate with Mobility

Cost_{switch}: Algorithm Switching Cost

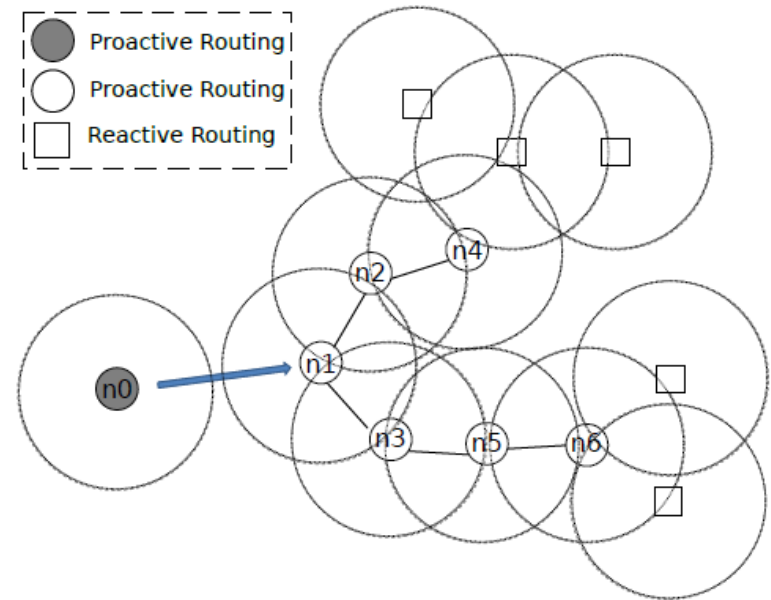
Objective: $\min(\text{Cost})$

$$\text{Cost} = \alpha \text{Cost}_{\text{switch}} + \beta \text{Cost}_{\text{Algorithm}}$$

$$\text{Cost}_{\text{switch}} = f_s^{\text{Switch}}(\text{Overhead}_{\text{Switch}}^{\text{CtrlMsg}})$$

$$\text{Cost}_{\text{Algorithm}} = f_s^{\text{Algorithm}}(\text{Overhead}_{\text{Algorithm}}^{\text{CtrlMsg}}(t))$$

- Node joins or leaves a proactive sub-network
 - Broadcast through the sub-network
- Switching cost is determined by the proactive routing algorithm



Conclusions and Future Work

- Conclusions
 - Propose self adaptive routing framework
 - Adapt an analytical model based on cost functions
 - Provide detailed expressions for $Cost_{Algorithm}$
 - Estimate depending on protocol and scenario
- Future Work
 - Determine the scaling factors o_r and o_p analytically
 - Adapt the parameters to typical fire fighting scenarios

Question?