

# Optimization and Performance Analysis of High Speed Mobile Access Networks

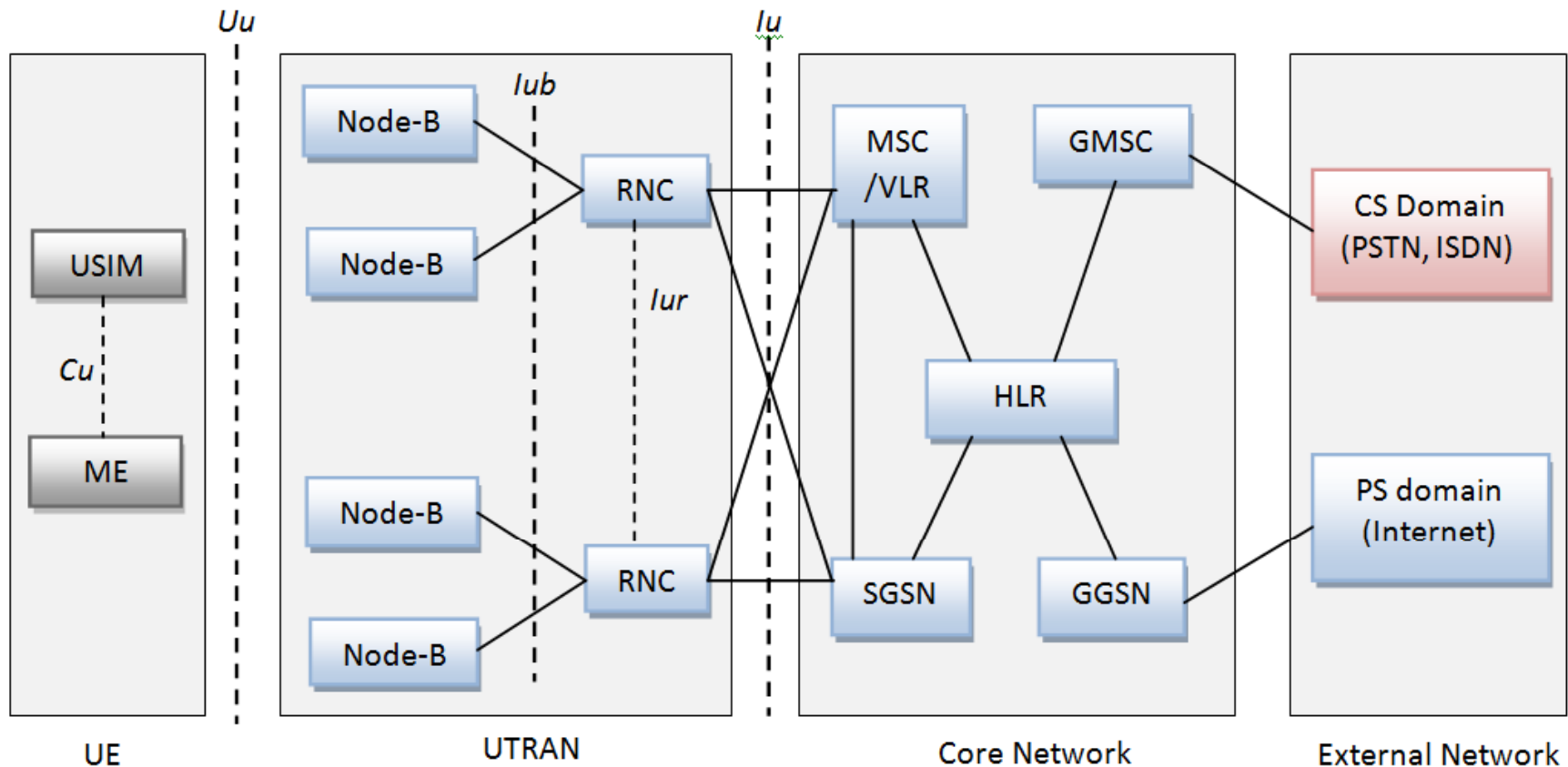
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ComNets, University of Bremen, Germany

03.09.2010

- ▶ Overview of high speed broadband wireless networks
  - Key technologies and architecture
- ▶ Main achievements
  - Overview of the completed tasks during the thesis work
- ▶ HSPA transport flow control and congestion control
  - Theoretical approach
  - Analytical modeling
  - Performance analysis and results comparison
- ▶ Conclusion and outlook

- ▶ Universal Mobile Telecommunication System (UMTS)
  - Standardised in 3GPP release 99 and radio interface: WCDMA
  - 64 kbit/sec circuit switched, 384 kbit/sec packet switched services
  - Bearer services, Location services and compatible with GSM
- ▶ High Speed Packet Access Networks (HSPA)
  - Downlink: HSDPA, standardised in 3GPP release 5
    - Higher data rates for packet services: 1.8, 3.6, 7.2 and 14 Mbit/sec
    - Key features: Hybrid ARQ, fast packet scheduling, adaptive modulation and coding (AMC)
    - IP Multimedia System (IMS)
  - Uplink: HSUPA, standardised in 3GPP release 6
    - Enhanced uplink data rates up to 5.76 Mbit/sec and
    - Key improvements as in the downlink
- ▶ Long Term Evolution (LTE)
  - Standardised in 3GPP release 8:
  - all IP-network, New OFDMA, MIMO based radio interface
  - Not backward compatible with previous UMTS

# TZi HSPA architecture (UMTS, HSDPA, HSUPA)



UE: User Equipment

VLR: Visitor Location Register

CS: Circuit Switch

SGSN: Serving GPRS Support Node

USIM: Universal Subscriber Identity Module

MCS: Mobile service Switching Centre

GGSN: Gateway GPRS support Node

PSTN: Public Switched Telephone Network

ME: Mobile Equipment

HLR: Home Location Register

EIR: Equipment Identity Register



UTRAN: UMTS Terrestrial Radio Access Network

RNC: Radio Network Controller

PS: Packet Switch

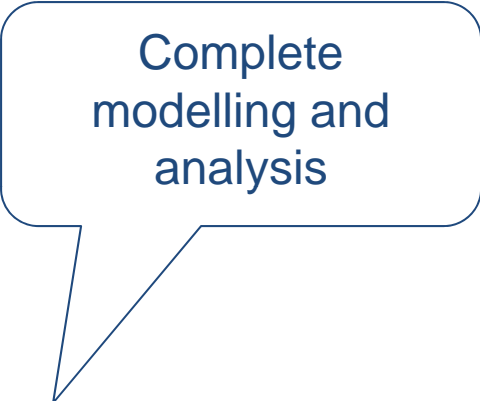
GMSC: Gateway MSC

# TZi Key achievements

- ▶ Design and development of a comprehensive HSPA simulator
  - Implementation of all UTRAN based protocols and end-user protocols
  - Design and implementation of uplink and downlink MAC schedulers
- ▶ HSPA transport feature development
  - Adaptive credit-based flow control schemes  *Journal of Communications 2009 (IEEE), Academy Publisher*
  - Effective congestion control schemes  *Journal Publication, IEEE VTC magazine, December 2009*
- ▶ Design and development of analytical models
  - A Markov model of congestion control
  - A combined Markov model of flow control and congestion control
- ▶ Design and development of a detailed LTE system simulator
  - E-UTRAN and end-user protocols
  - A comprehensive MAC scheduler and IP based transport QOS scheduler

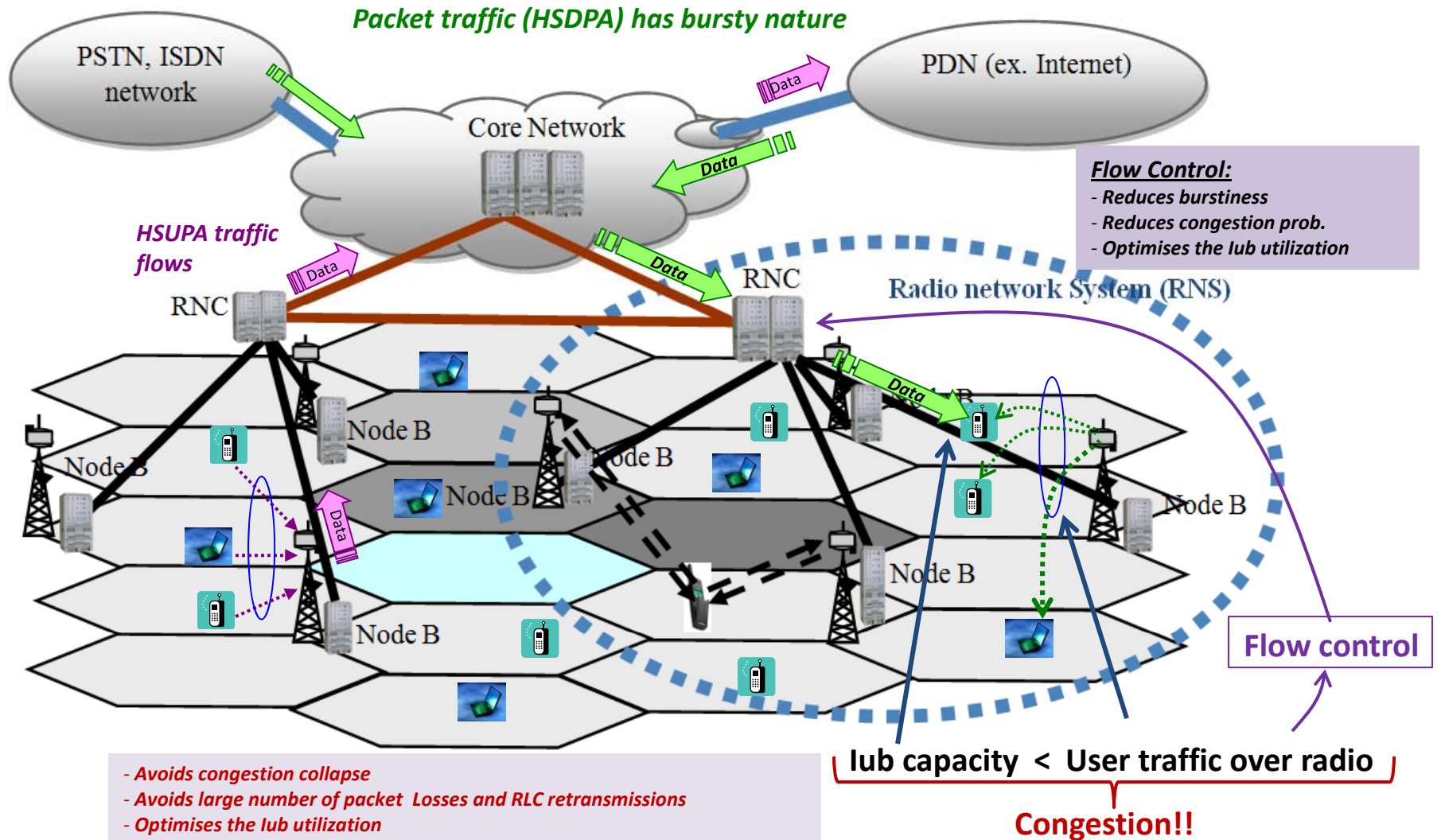
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- ▶ Design and development of a comprehensive HSPA simulator
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  - Design of uplink and downlink MAC schedulers
- ▶ HSPA transport feature development
  - A adaptive credit-based flow control scheme and
  - Effective congestion control schemes
- ▶ Design and development of analytical models
  - A Markov model of congestion control
  - **A combined Markov model of flow control and congestion control**
- ▶ Design and development of a detailed LTE system simulator
  - E-UTRAN and end-user protocols
  - A comprehensive MAC scheduler and IP based transport QoS scheduler



Complete  
modelling and  
analysis

# TZi HSDPA FC and CC overview



## ▶ Adaptive credit-based flow control

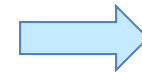
- Flow control adapts the lub flow to the available throughput at the air interface for individual user flows

- Credit-based flow control

- Continuous loop control using the **Provided Bit Rate (PBR)**.

$$\overline{PBR}(t) = w \cdot \overline{PBR}(t-1) + (1-w) \cdot PBR(t) \quad \text{where } w \text{ is the weight factor}$$

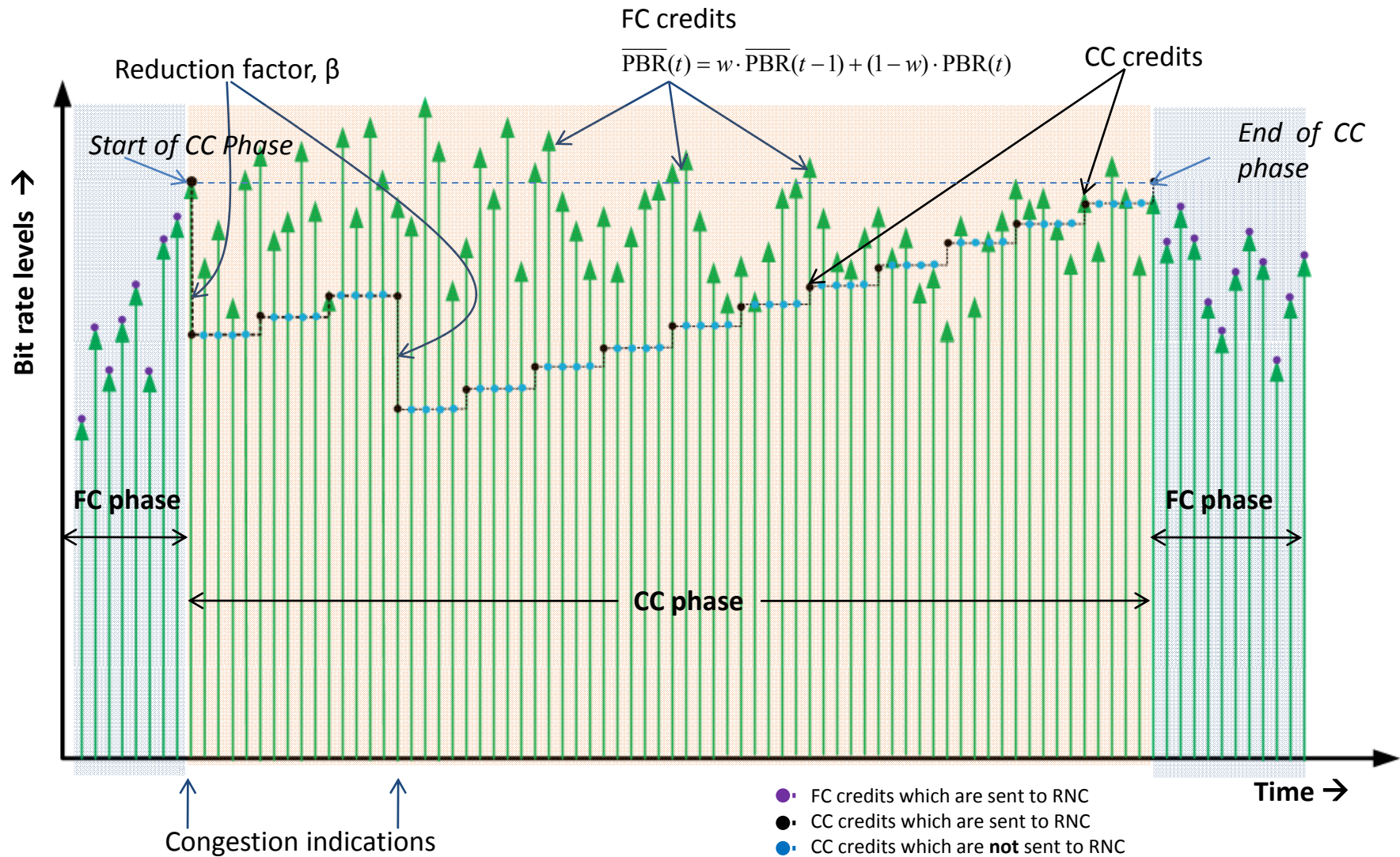
- Buffer management (to optimise the radio utilisation)



## ▶ Congestion control

- Bursty traffic over the limited transport network causes congestion resulting in many packet losses
- Wastage of scarce radio resources, network resources and degradation of overall end-to-end performance
- Requirement: a proper congestion control mechanism to adapt radio capacity to the available transport capacity adaptively
- Congestion control mechanism includes
  - **Preventive** and **reactive** congestion **detection** schemes and input traffic **control** scheme





## ▶ Prerequisites

- Two state variables for FC and CC
- Time step in CC is several times longer (5) than the time step in FC
- Maximum level reached under CC depends on starting FC level

## ▶ Assumptions

- The interarrival times of CIs are independent and identically distributed
- Number of users remains constant (stationary system)
- Constant transmission delay for CA signals
- Per-user buffer occupancy at Node-B is not considered for FC modelling

# TZi Joint Markov chain

## State representation

- three non-negative integers,  $[i, j, k]$
- Bit rate level in FC state,  $i$  [ $i = 1, 2, 3, \dots, m$ ]
  - Bit rate level in CC state,  $j$  [ $j = 1, 2, 3, \dots, m$ ]
  - Time steps in CC state,  $k$  [ $k = 1, 2, 3, \dots, 5$ ]

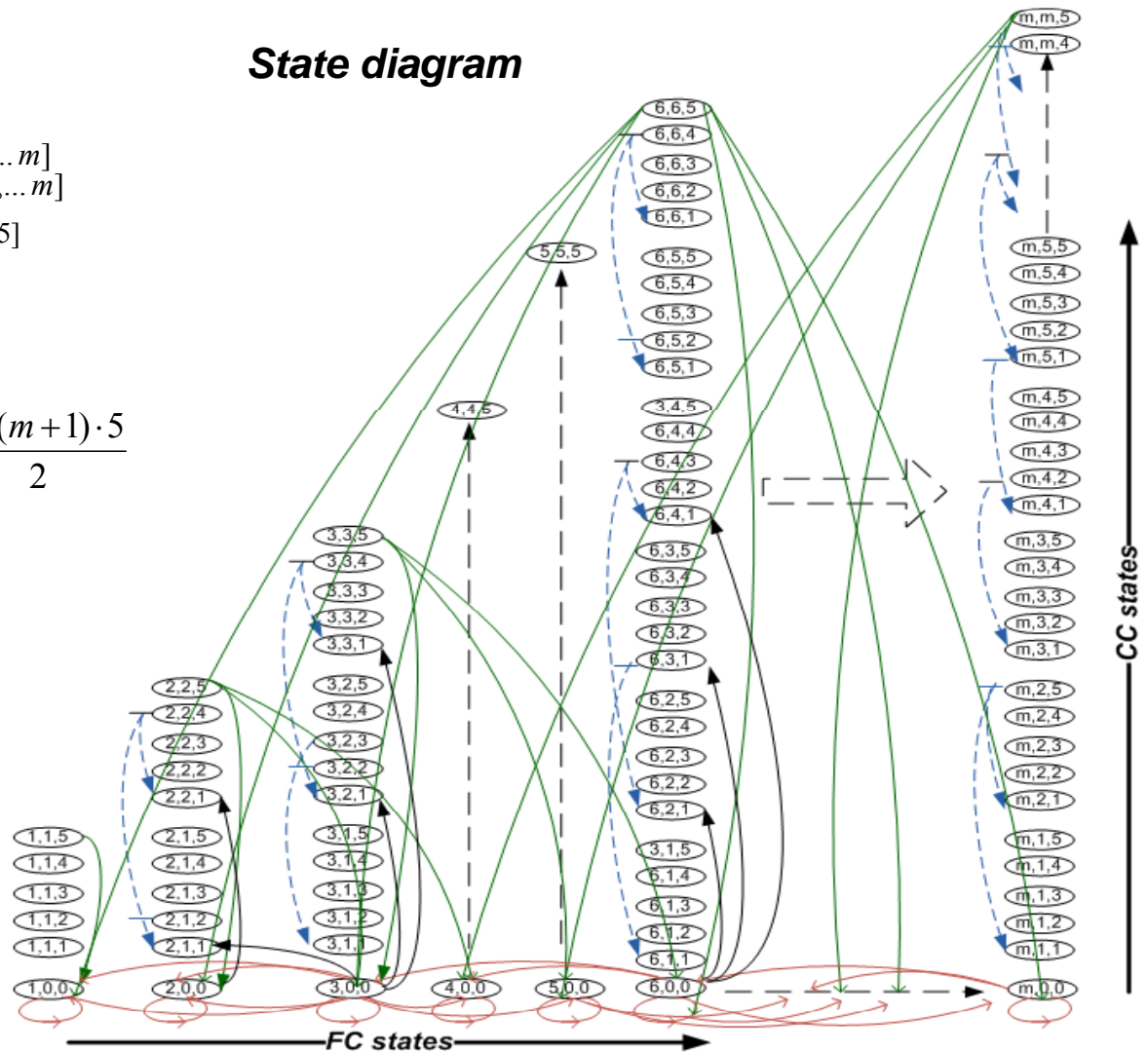
## The number of states

$$\left. \begin{array}{l} \text{The total number of FC} \\ \text{states and CC states} \end{array} \right\} = m + \frac{m \cdot (m+1) \cdot 5}{2}$$

## State transition

- FC to FC transitions  $\rightarrow$  (red arrow)
- FC to CC transitions  $\rightarrow$  (black arrow)
- CC to CC transitions  $\rightarrow$  (blue arrow)
- CC to FC transitions  $\rightarrow$  (green arrow)

## State diagram



- ▶ Stationary FC state probability matrix,  $PBRm$

$$PBRm = [pbr_j]_{1 \times m} \quad \text{where } j = 1 \dots m$$

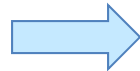
(The stationary FC state probability matrix is derived from a trace file which is taken from a dedicated radio simulation.)

- ▶ The congestion indication arrival probability matrix,  $A_{ci}$

$$A_{ci} = [q_i]_{1 \times (1+d_{\max})}$$

$$\text{where } q_i = \Pr[\text{exactly } i \text{ CI signals during } \Delta T]; \quad i = 0 \dots d_{\max}$$

(This parameters, the CI arrival probability within a step is taken from a trace which is taken from a **fast queuing simulator** which is designed and developed by the author)



# TZi Transition probability calculations

- ▶ Flow control transition probability matrix

$$P_{fctofc} = [p_{ij}]_{m \times m}$$

$$p_{ij} = \Pr[FC(T + \Delta T) = j \mid FC(T) = i]$$

- ▶ State transition probability from FC state,  $i$  to FC state,  $j$

$$P_{(i,0,0)(j,0,0)} = q_0 p_{ij} \quad \text{where } q_0 \text{ is the probability that no CI arrivals occurs within a given FC interval}$$

- ▶ State transition probability from FC state to CC state

$$P_{(i,0,0)(i,l,1)}^+ = q_n \quad \forall l = i \times \alpha^n; \quad n = 1 \dots d_{\max}$$

where  $\alpha = 1 - \beta$  and

$d_{\max}$  max number of CI arrivals within a single step



▶ From CC state to CC state transition probabilities

- Up: in case of no arrivals

$$P_{(i,j,k)(i,j,k+1)} = q_0 \quad \text{for } k = 1,2,3,4$$

$$P_{(i,j,k)(i,j+1,1)} = q_0 \quad \text{for } k = 5$$

- Down: due CI arrivals

$$P_{(i,j,k)(i,l,1)}^+ = q_n \quad \forall l = j \times \alpha^n; \quad n = 1 \dots d_{\max} \quad \text{and } k = 1 \dots 5$$

▶ From CC state to FC state transition probabilities

$$P_{(i,j,5)(l,0,0)} = q_0 p_{il} \quad \begin{array}{l} l \text{ is the next FC state and } i \text{ is the starting FC state} \\ \text{before the MAC - d flow enters the CC state} \end{array}$$

- ▶ Transition probability matrix, with square dimension  $n_t$

$$P = [p_{ijk}]_{n_t \times n_t} \quad \text{where } i = 1, 2, 3, \dots, n_{fc}$$

$$j = 1, 2, 3, \dots, n_{cc}$$

$$k = 1, 2, 3, \dots, n_{st}$$

- ▶ Stationary state probabilities matrix,  $\pi$

$$\pi = \pi \cdot P \quad \text{Where } \pi \text{ denotes the state vector, } [\pi_0, \pi_1, \pi_2, \pi_3, \dots, \pi_{n_t}]$$

- ▶ Average throughput =  $\text{bitRateStepSize} \times \sum_{i=1}^{n_t} i \pi_i \quad \text{bit / sec}$

- Example, size of the bit rate level is 33.6 kbps for the given consideration

## ▶ Parameter configuration

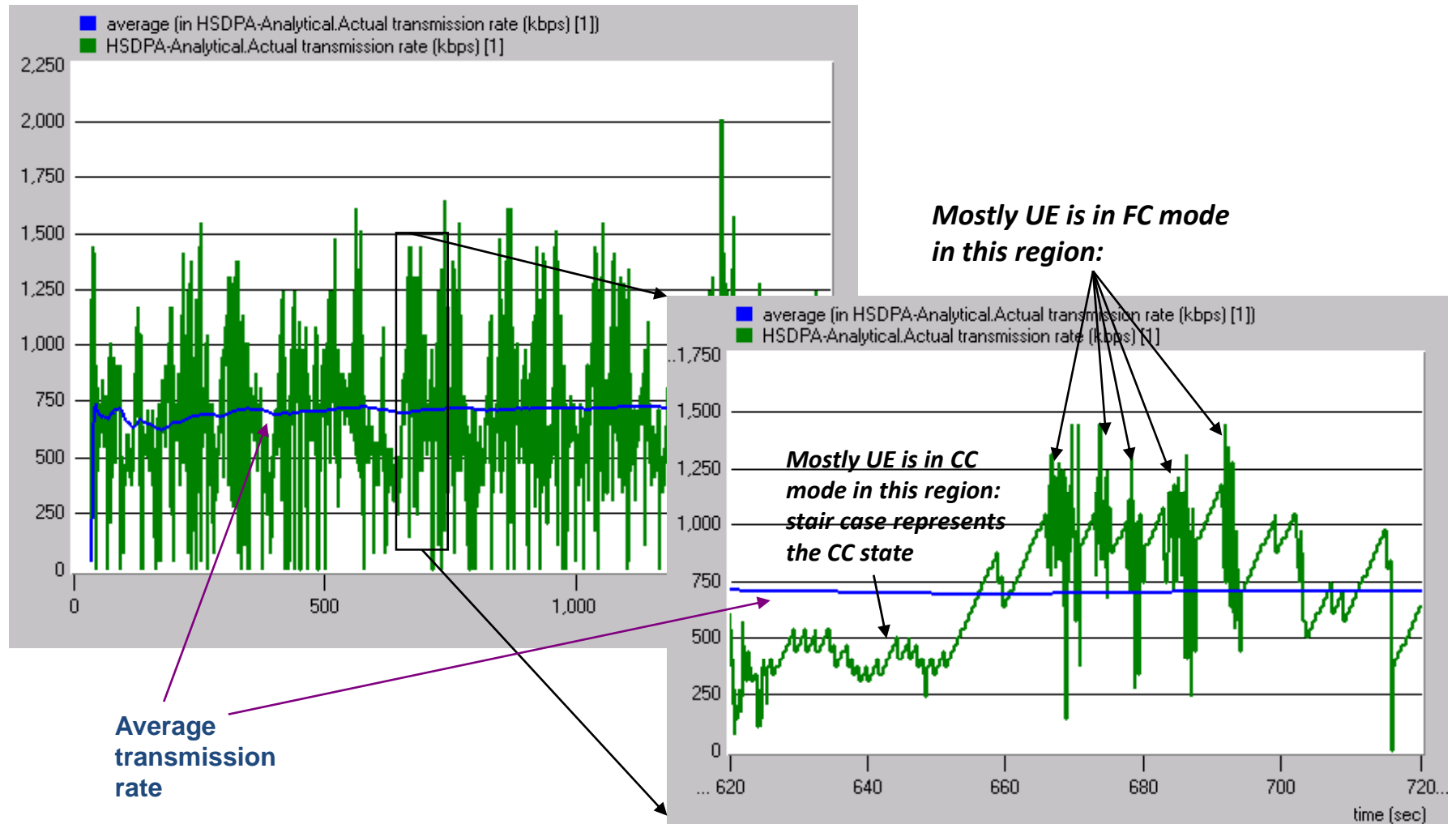
### ■ Common parameters (simulation and analytical)

- FC cycle time = 100 ms
- CC AIMD cycle time = 500 ms
- FC and CC Step size = 33.6 kbps
- Reduction factor  $\beta = 0.25$
- Safe timer = 80 ms
- ATM bandwidth = 2Mbps

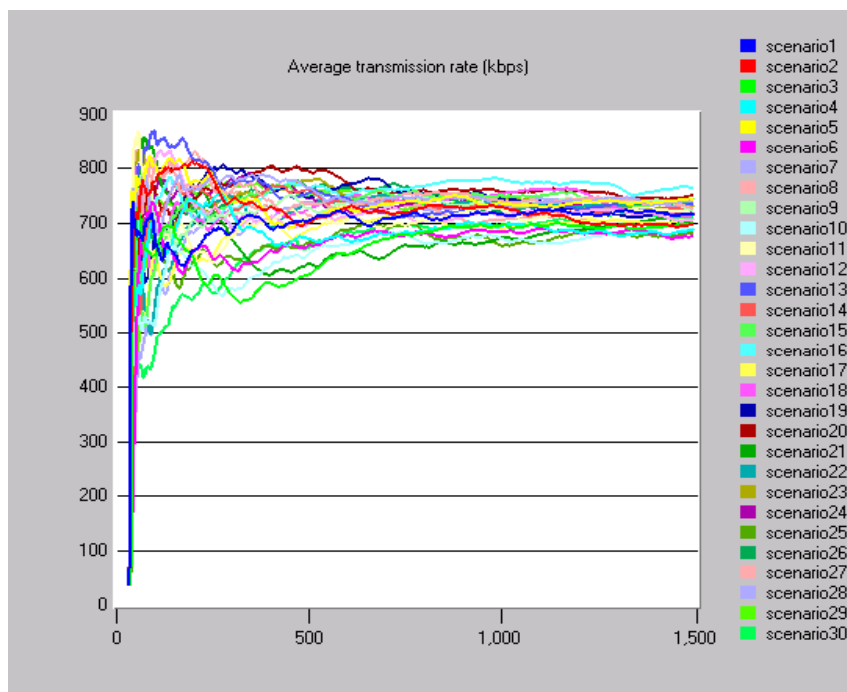
### ■ Simulation specific

- Traffic models: FTP and HTTP traffic models
- User constellation: 18 users, 1 FTP user who downloads a large file during the simulation time and uses the probability distribution used for the analytical model. All other 17 users generate HTTP traffic.
- The simulation duration is 2000 sec and 32 replications are used to determine the confidence interval.

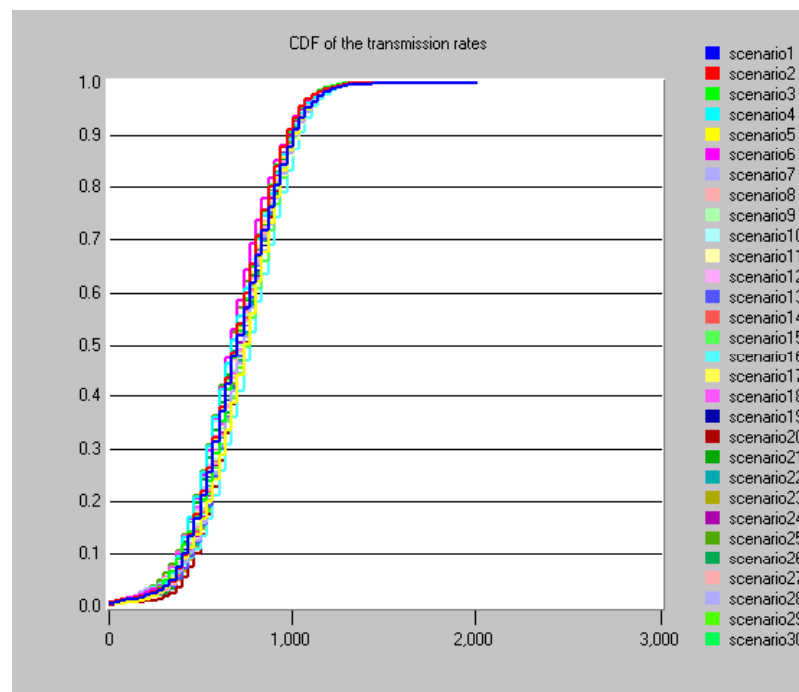




### Average throughputs



### CDF of throughputs



### Statistical evaluation

Mean = 719.60 kbit/sec

Standard deviation = 21.13 kbit/sec

95%confidence interval

= [713.05 kbit/sec– 726.16 kbit/sec]

<i>Replications, i</i>	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Mean, $\bar{Y}_i$	713	697	700	688	743	676	740	726	720	689	723	706	734	726	743
<i>Replications, i</i>	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30
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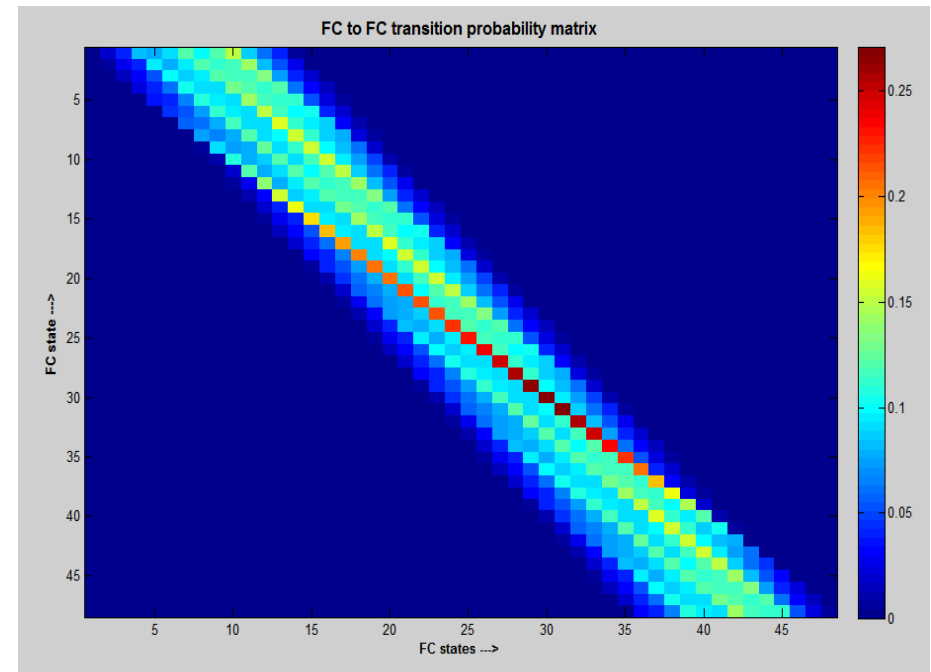
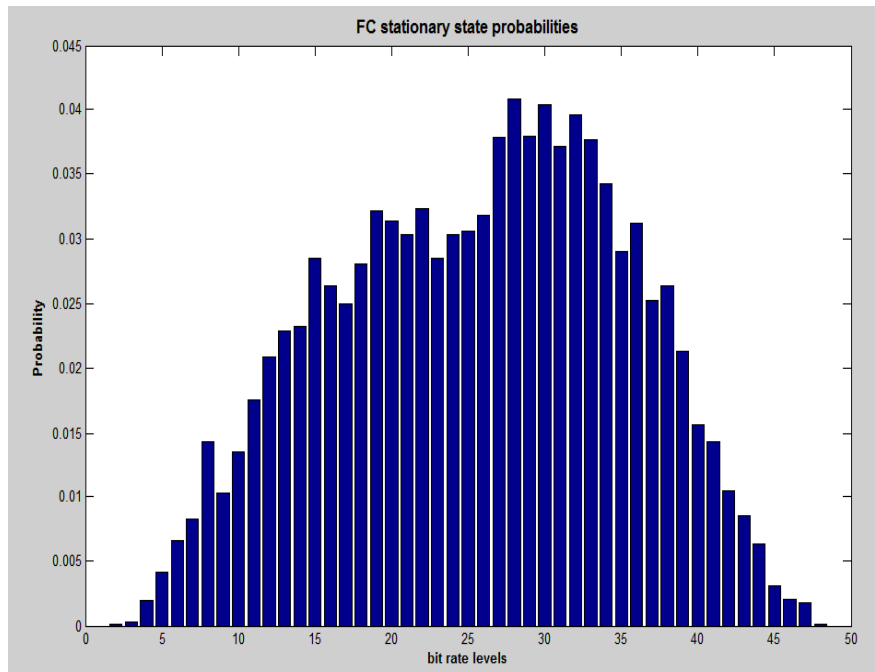
Probability mass function of the bit rate levels (number of MAC-ds)

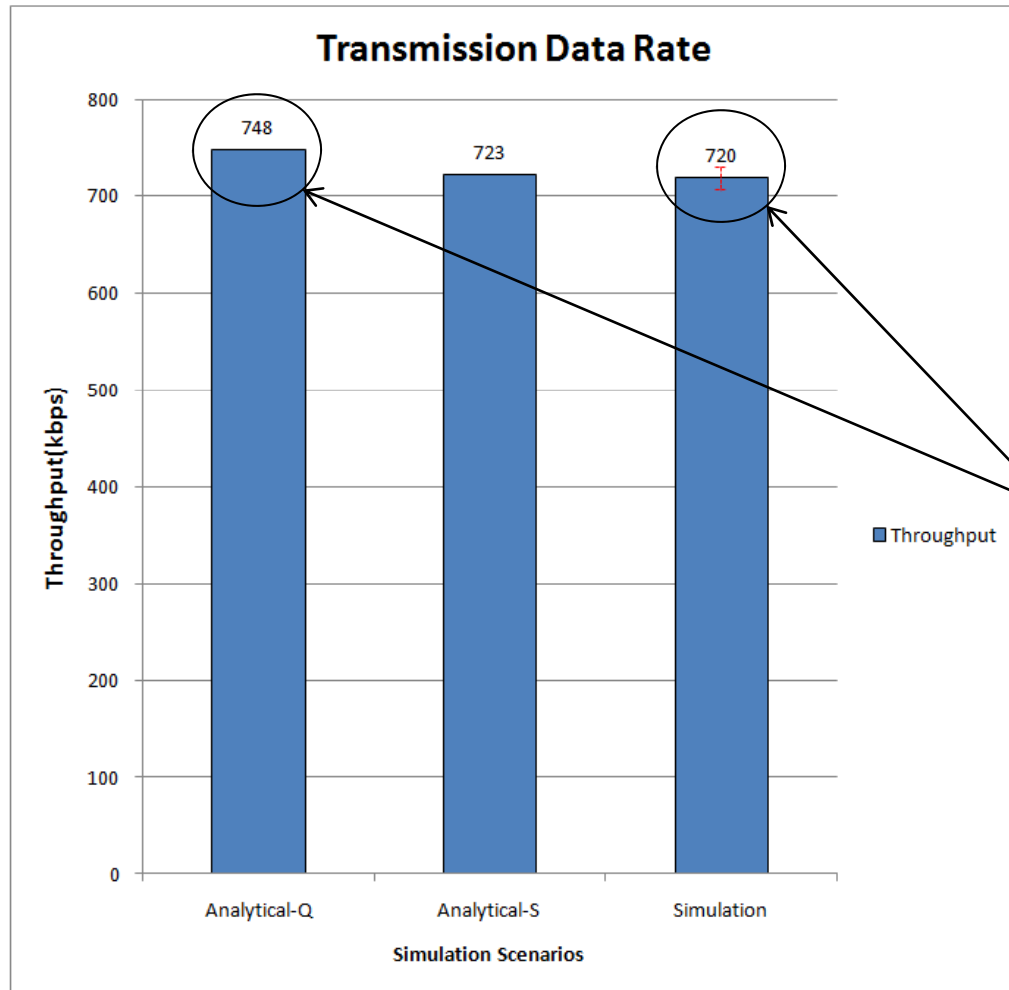
*FC state stationary probabilities*

FC filter formula

$$\overline{PBR}(t) = w \cdot \overline{PBR}(t-1) + (1-w) \cdot PBR(t)$$

$w$  is the weight factor





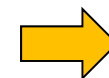
## Scenarios

- Analytical model-Q:  
CI trace from the fast queuing simulator
- Analytical model-S  
CI trace from the detailed system simulator
- Simulation

	Analytical-S	Analytical-Q
Probability of no CI arrival within a single step, $q_0$	95.83%	94.74%

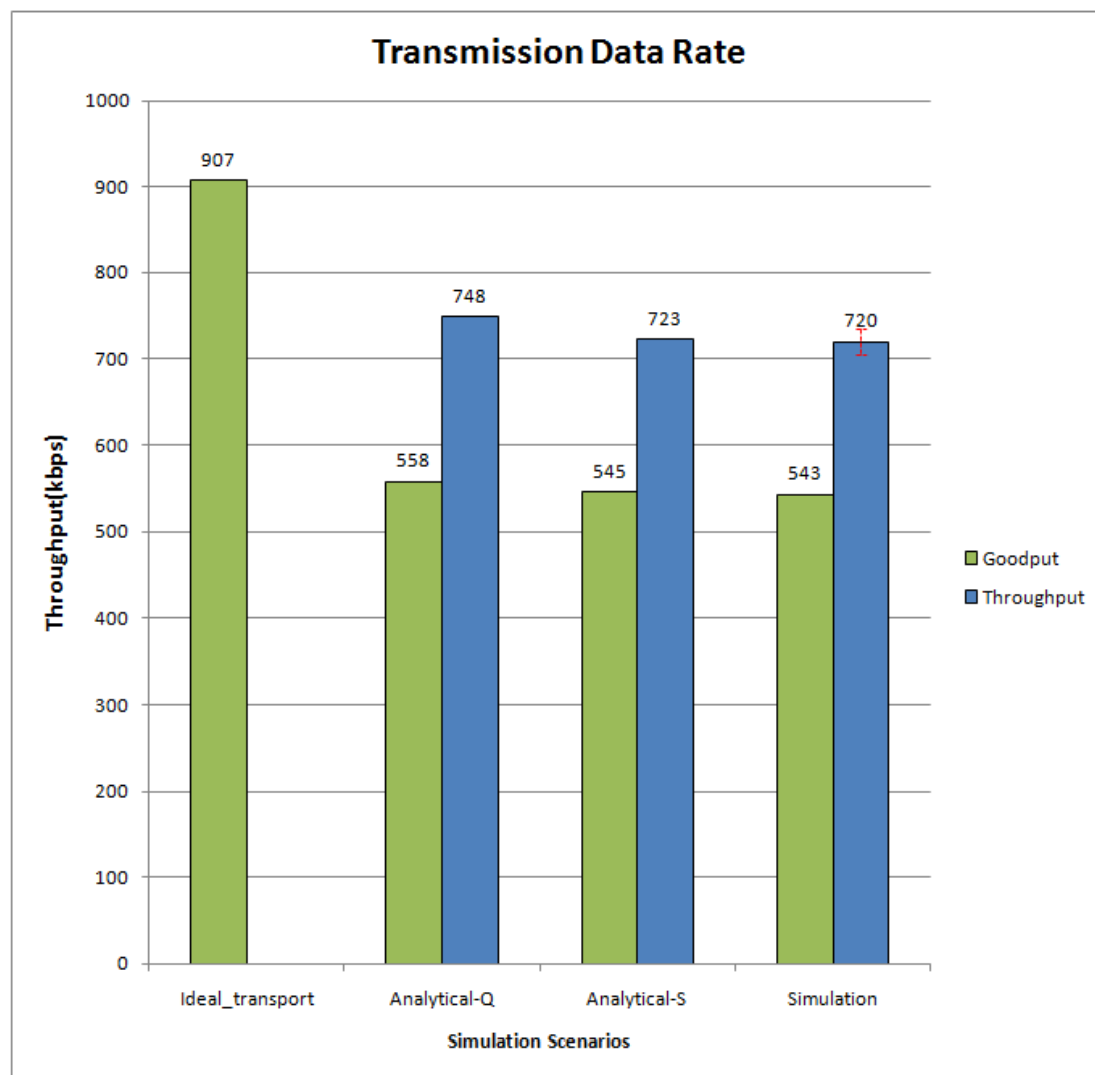
## Throughput difference (analytical-Q)

- Additional RLC retransmissions (1%, - 0.41%)
- Losses are uniformly distributed
- Effect of TCP the TCP protocol
- All other protocol simplifications and assumption ..



Goodput comparison provide more close agreement

# TZi Results comparison cont.



	Throughput	Goodput
<b>Simulation Results</b>	719.65	543.09
<b>Analytical Results-S</b>	722.63	545.38
<b>Analytical Results-Q</b>	748.48	558.10
<b>ideal Transport</b>	-	907.18

## Key achievements

- ▶ A good match between the analytical model results and the detailed system simulation results
- ▶ The FC and CC investigation and analysis can be done **faster** using the analytical model along with the fast queuing simulator compared to detailed system simulator
  - Detailed system simulator: **order of days**
  - Analytical model: **order of minutes**

- ▶ Detailed HSPA system simulator has been implemented, tested, validated and used for the performance evaluation.
- ▶ TNL credit-based adaptive flow control and congestion algorithms have been implemented tested, validated and used for end-to-end performance analysis.
  - Overall network performance can be significantly improved by **reducing burstiness** over the transport network, **optimising transport utilisation** and effectively **minimising congestion** in the transport network
  - FC and CC algorithms provides **guaranteed end-user QoS** while achieving a **optimum end-user performance**
- ▶ Two analytical models has been implemented, tested, validated and evaluated the performance.
  - A **Markov model** for modelling **congestion control** functionality
  - A **joint Markov model** for modelling **FC** and **CC** functionalities

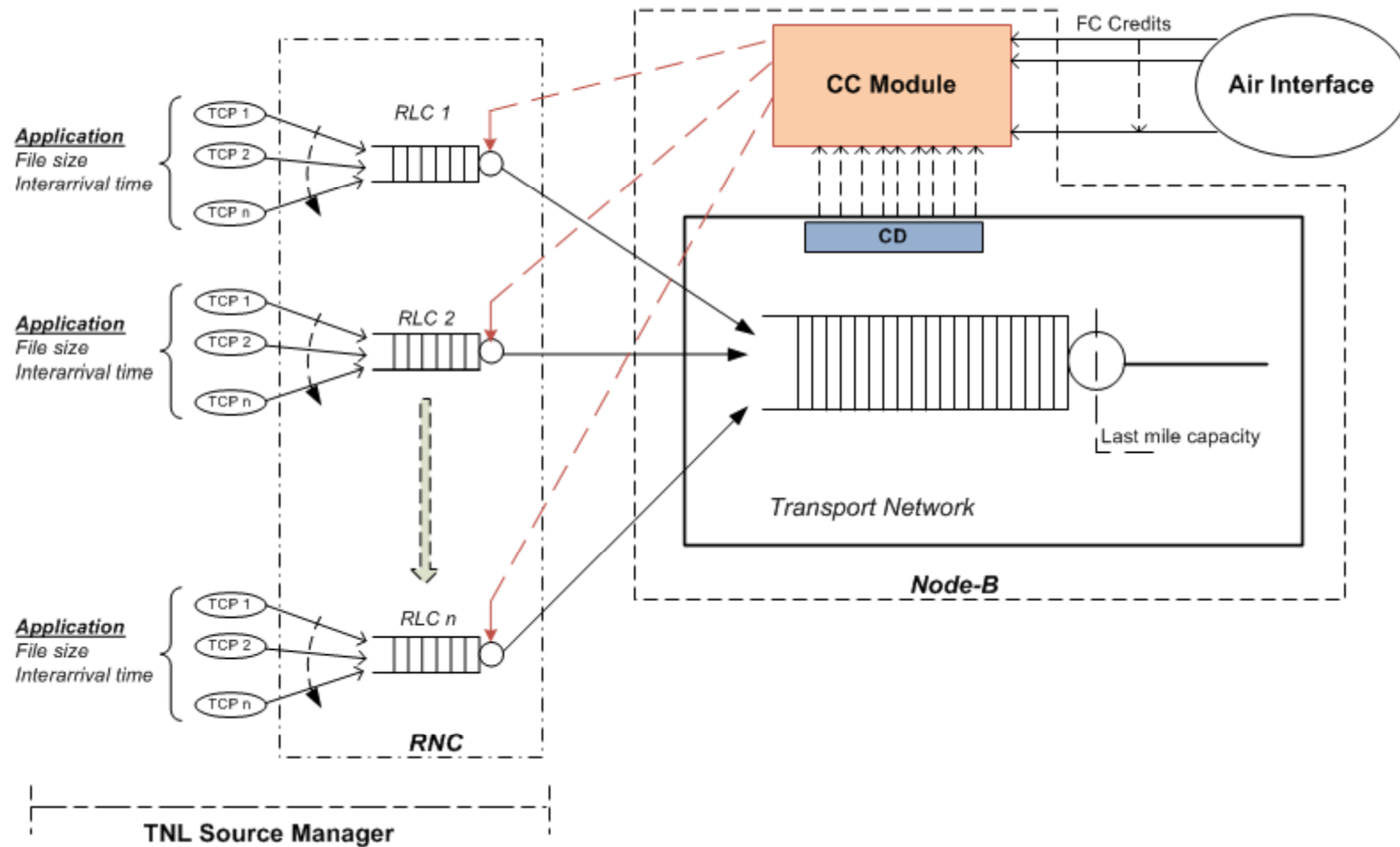
- ▶ There is a good match between analytical model results and the detailed system simulation results.
- ▶ A complete faster alternative solution to the timing consuming detailed system simulator can be provided by the analytical mode along with the fast queuing for TNL feature analysis.
- ▶ Detailed LTE system simulator implemented, tested, validated and performance analysed.
  - ▶ In addition to general protocol implementation, MAC scheduler and transport QoS packet scheduler have been implemented
  - ▶ Effects of transport congestion for network and end-user performance have been studied and analysed

- ▶ Proper flow control and congestion control schemes are needed to be proposed and implemented in the LTE UTRAN in order to protect transport packet losses due to congestion
  - ▶ UL congestion control and load balancing for non-GBR bearers
  - ▶ DL congestion handling mainly for non-GBR bearers
  - ▶ Effective admission control mechanism for GBR bearers
- ▶ There is a clear requirement of cross layer functionalities between Radio MAC scheduler and transport scheduler for effective QoS management

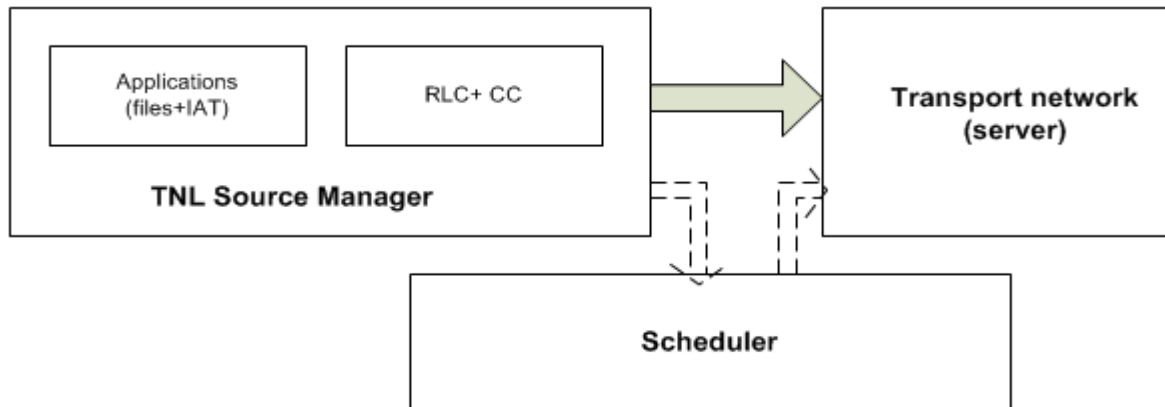


**Thank you very much!**

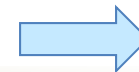
Any Questions?

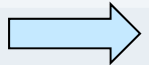


# TZi Fast queuing simulator

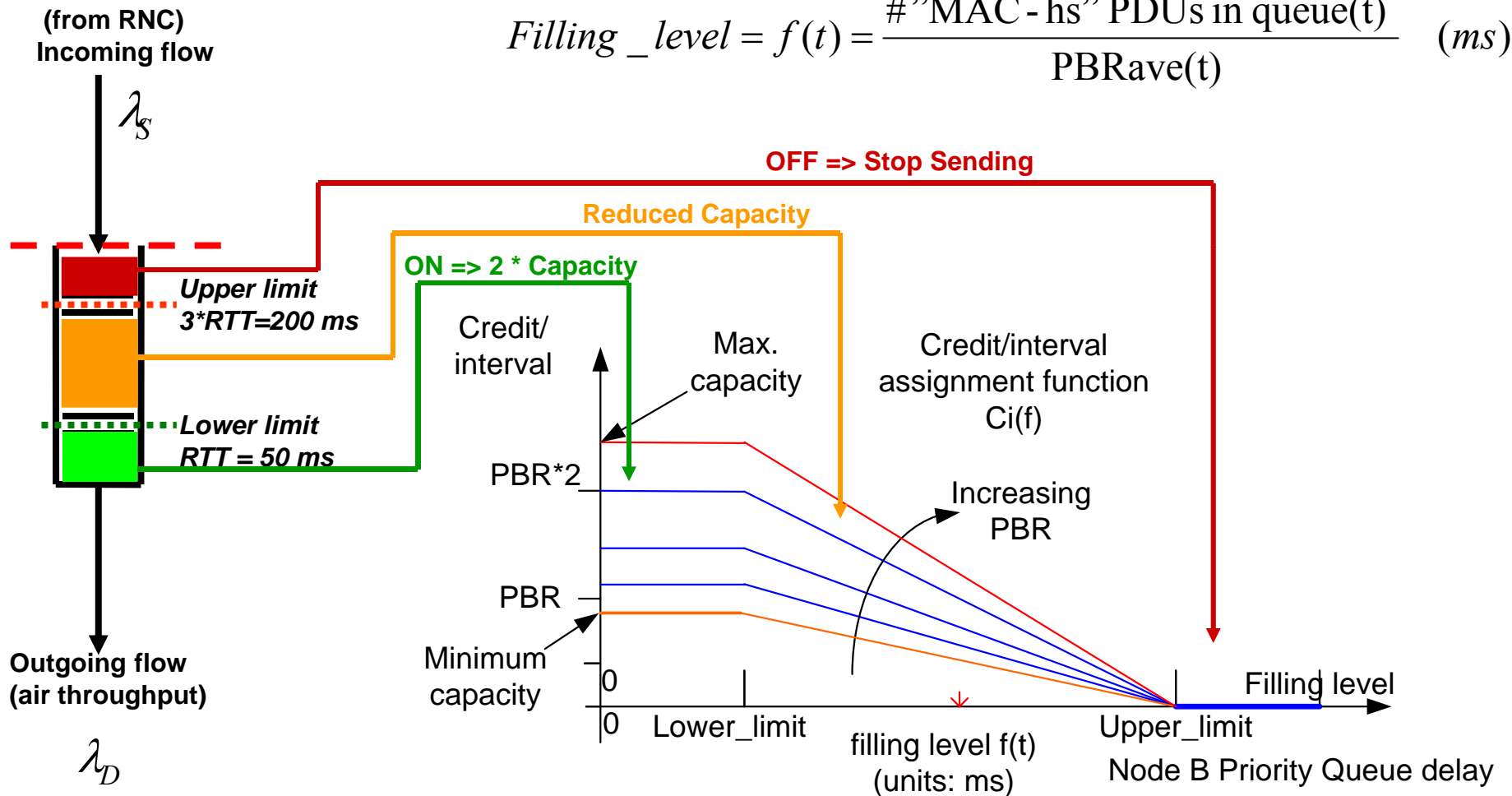


- ▶ A fast queuing simulator is implementation using CNCL library
- ▶ Key assumptions and simplifications
  - Traffic sources are modelled without complex TCP protocol functionality
  - Transport loss ratio set to the maximum value 1% and losses are uniformly distributed. (RLC able to recover the loss before TCP notices)
  - The CI arrival process assumed as a Poisson arrival

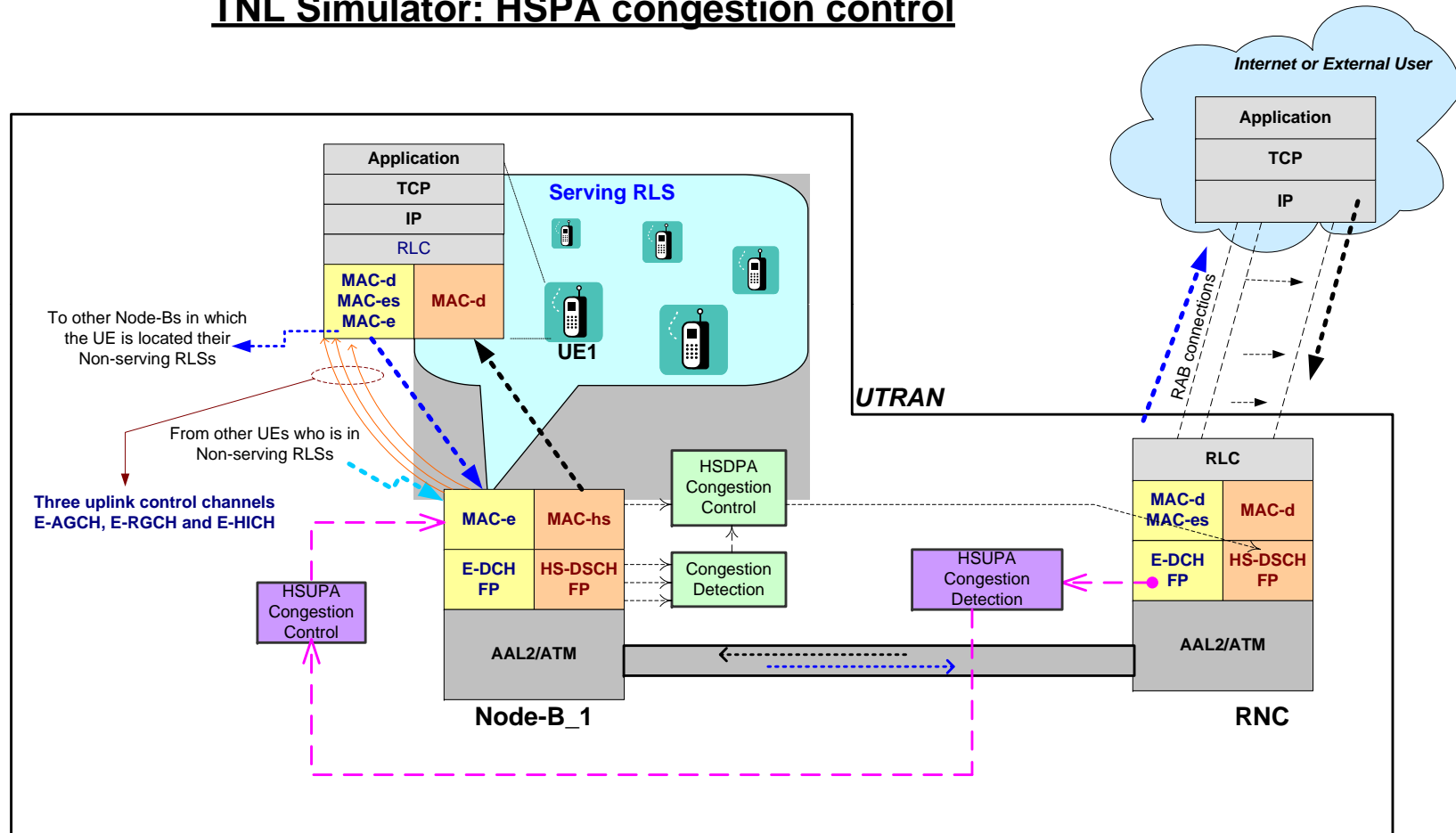




$$Filling\_level = f(t) = \frac{\# \text{ "MAC - hs" PDU's in queue}(t)}{PBRave(t)} \quad (ms)$$



## TNL Simulator: HSPA congestion control

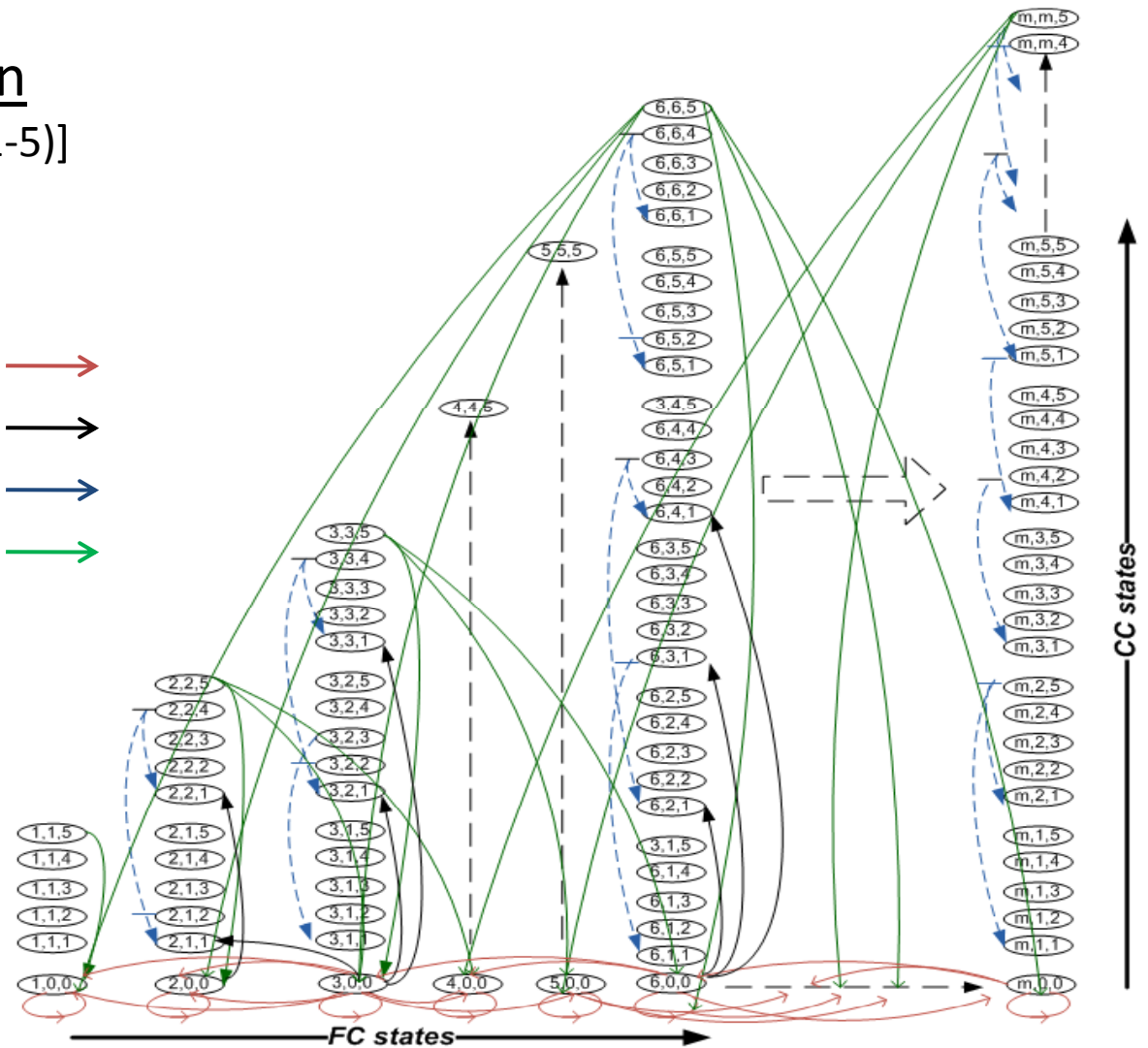


## State representation

$[i,j,k] = [FC, CC, \text{Steps}(1-5)]$

## State transition

- FC to FC transitions →
- FC to CC transitions →
- CC to CC transitions →
- CC to FC transitions →



# TZi Markov model: states

- ▶ Each state is identified by three non-negative integers,  $[i, j, k]$ 
  - Bit rate level in the FC state,  $i$  [ $i = 1, 2, 3, \dots, m$ ]
  - Bit rate level in the CC state,  $j$  [ $j = 1, 2, 3, \dots, m$ ]
    - “ $j=0$ ” indicates the state within FC state
  - Time step in CC state,  $k$  [ $k = 1, 2, 3, \dots, 5$ ]
    - Again “ $k=0$ ” indicates the state within FC state
  
- ▶ Number of states in Markov model
  - Total number of FC states =  $m$
  - Total number of CC states =  $\frac{m \cdot (m + 1) \cdot 5}{2}$ , where CC state finite sequences  $5, 10, 15, 20, \dots, 5m \Rightarrow (5n)_{n=1}^m$
  - The total number of FC states and CC states:  $m + \frac{m \cdot (m + 1) \cdot 5}{2}$

*For example, if  $m$  is 48 ( $\text{eff\_BW/step} = 2.0 \cdot 10^6 \cdot 0.8 / 33.6 \cdot 10^3$ ), then the total number of states is 5928.*

# TZi Simulation results analysis

## ▶ Replication mean calculation

Replication 1 :  $Y_1, Y_2, Y_3, \dots, Y_m \text{ ----- } > \bar{Y}_1$

Replication 2 :  $Y_{m+1}, Y_{m+2}, Y_{m+3}, \dots, Y_{2m} \text{ ----- } > \bar{Y}_2$

Replication 3 :  $Y_{2m+1}, Y_{2m+2}, Y_{2m+3}, \dots, Y_{3m} \text{ ----- } > \bar{Y}_3$

..

Replication n :  $Y_{(n-1)m+1}, Y_{(n-1)m+2}, Y_{(n-1)m+3}, \dots, Y_{nm} \text{ ----- } > \bar{Y}_n$

$$\text{Mean : } \bar{Y} = \frac{1}{n} \sum_{i=1}^n \bar{Y}_i$$

$$\text{Variance: } V = \frac{1}{n-1} \sum_{i=1}^n (\bar{Y} - \bar{Y}_i)^2$$

$$\text{Confidence interval} = \bar{Y} \pm (t_{\alpha/2, n-1} \times \sqrt{\frac{V}{n}})$$

where 100(1- $\alpha$ )% CI

where  $t_{\alpha/2, n-1}$  is taken from Student t-distribution. “n-1” is the degree of freedom and 1- $\alpha$  is the confidence level.

## ▶ Statistical evaluation

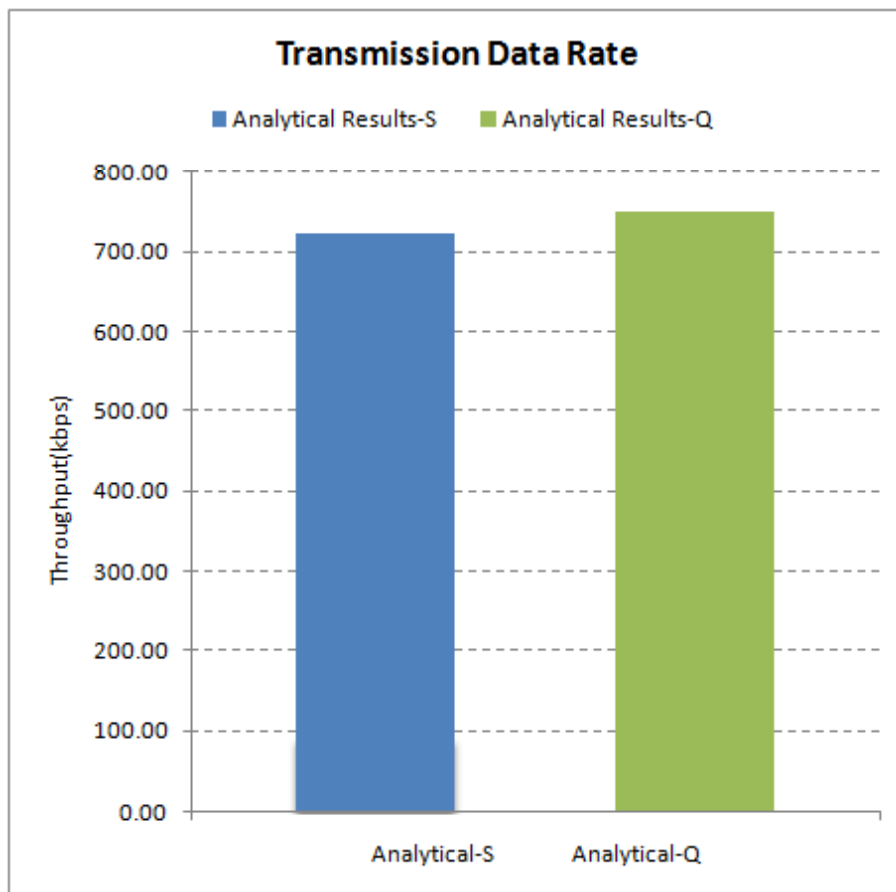
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- Mean = 719.60 kbit/sec
- Standard deviation = 21.13 kbit/sec
- 95%confidence interval = [713.05 kbit/sec– 726.16 kbit/sec]

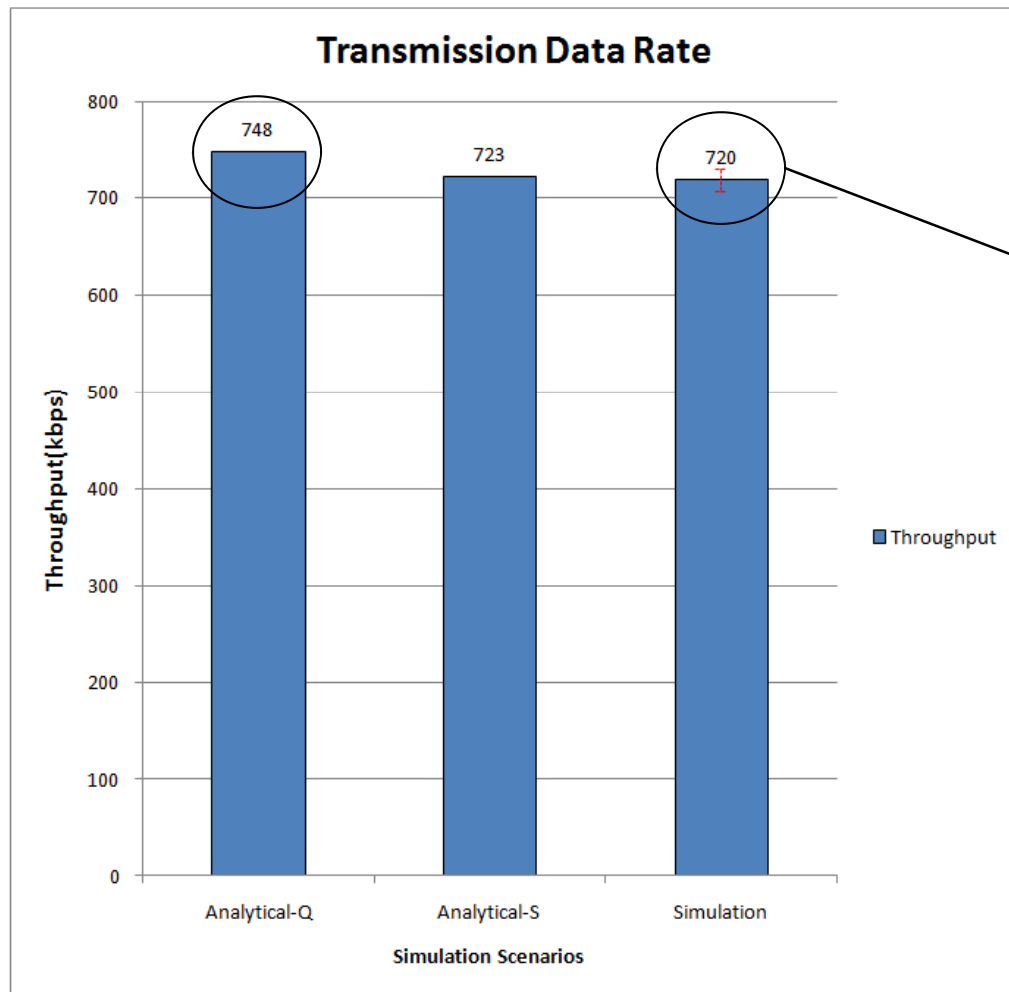


- ▶ Analytical model-Q
  - CI trace from the fast queuing simulator
- ▶ Analytical model-S
  - CI trace from the detailed system simulator

	Analytical-S	Analytical-Q
Probability of no CI arrival within a single step, $q_0$	95.83%	94.74%



# TZi Results comparison

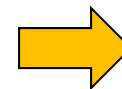


## Simulation

Mean throughput = 719.65 kbit/sec  
95% CI (713.05 kbps, 726.16 kbit/sec)

*There are many issues for this difference during this analysis*

- ▶ *Fast queuing simulator analysis had higher number of RLC retransmissions compared to detailed system simulator and it leads to a higher throughput for fast queuing simulator (loss probability of the detailed system simulator is 0.41% whereas it is 1% for the fast queuing simulator. Additional overhead is ~6kbit/sec.)*
- ▶ *There are other effects which cannot be quantified as above such as*
  - *Losses are uniformly distributed*
  - *Effect of TCP slow start etc..*
- ▶ *And all other protocol effects and assumptions*



Goodput comparison provide more close agreement

# TZ Joint Markov chain

## State representation

- three non-negative integers,  $[i, j, k]$
- Bit rate level in FC state,  $i$  [ $i = 1, 2, 3, \dots, m$ ]
  - Bit rate level in CC state,  $j$  [ $j = 1, 2, 3, \dots, m$ ]
  - Time steps in CC state,  $k$  [ $k = 1, 2, 3, \dots, 5$ ]

## The number of states

$$\left. \begin{array}{l} \text{The total number of FC} \\ \text{states and CC states} \end{array} \right\} = m + \frac{m \cdot (m+1) \cdot 5}{2}$$

## State transition

- FC to FC transitions  $\rightarrow$  (red arrow)
- FC to CC transitions  $\rightarrow$  (black arrow)
- CC to CC transitions  $\rightarrow$  (blue arrow)
- CC to FC transitions  $\rightarrow$  (green arrow)

## Input Parameters

Stationary FC state prob. matrix,  
 $PBR_m = [pbr_j]_{1 \times m}$  where  $j = 1 \dots m$

CI arrival prob. matrix,  
 $A_{ci} = [q_i]_{1 \times (1+d_{\max})}$  where  
 $q_i = \Pr[\text{exactly } i \text{ CI signals during } \Delta T];$   
 $i = 0 \dots d_{\max}$

## State diagram

